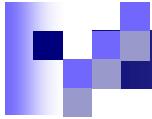


# The RuleML Family of Web Rule Languages

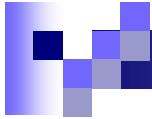
PPSWR'06, Budva, Montenegro, 10 June 2006

Harold Boley  
University of New Brunswick, Canada  
National Research Council of Canada



# Introduction

- Rules are central to the Semantic Web
- Rule interchange in an open format  
is important for e-Business
- RuleML is the de facto open language  
standard for rule interchange/markup
- Collaborating with W3C ([RIF](#)), OMG (PRR,  
SBVR), OASIS, DARPA-DAML, EU-REWERSE,  
and other standards/gov'nt bodies



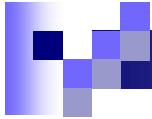
# RuleML Enables ...

## Rule

modelling  
markup  
translation  
interchange  
execution  
publication  
archiving

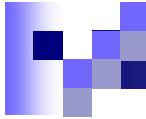
in

**UML**  
**RDF**  
**XML**  
**ASCII**



## RuleML Identifies ...

- Expressive sublanguages
  - for Web rules
  - started with
    - *Derivation* rules: extend SQL views
    - *Reaction* rules: extend SQL triggers
  - to empower their subcommunities



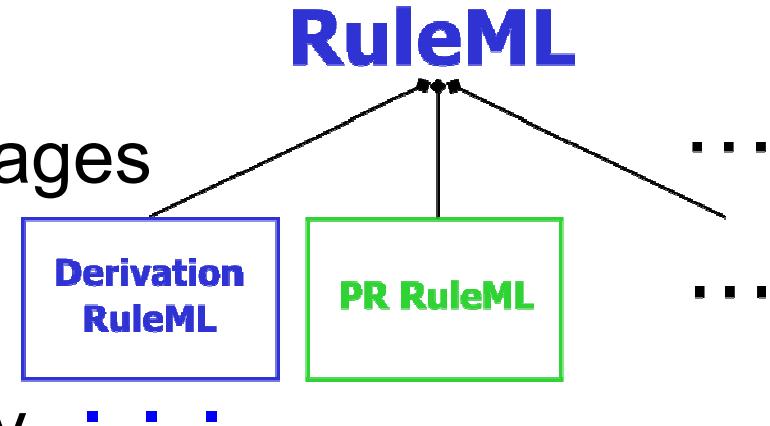
## RuleML Specifies ...

- Derivation rules via XML Schema:
  - All sublanguages:(OO) RuleML 0.9
  - First Order Logic: FOL RuleML 0.9
  - With Ontology language: SWRL 0.7
    - A Semantic Web Rule Language  
Combining OWL (W3C) and RuleML
  - With Web Services language: SWSL 0.9
- Translators in & out (e.g. Jess) via XSLT

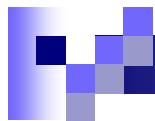


# Modular Schemas

“RuleML is a **family** of sublanguages whose **root** allows access to the language as a whole and whose **members** allow to identify customized subsets of the language.”

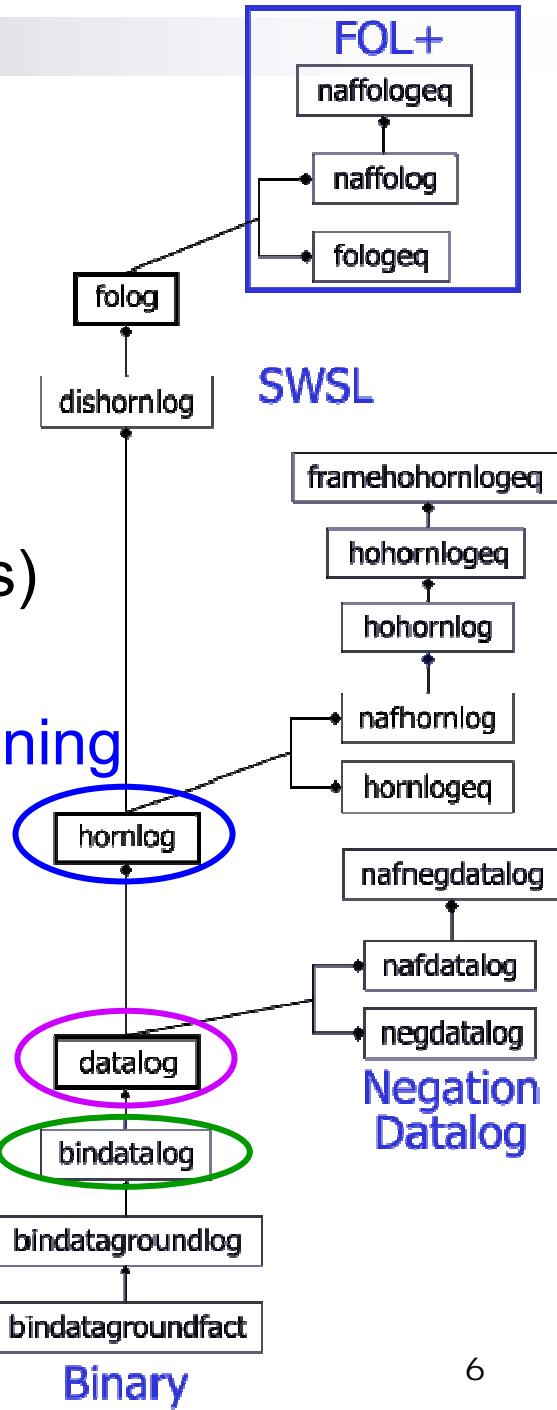


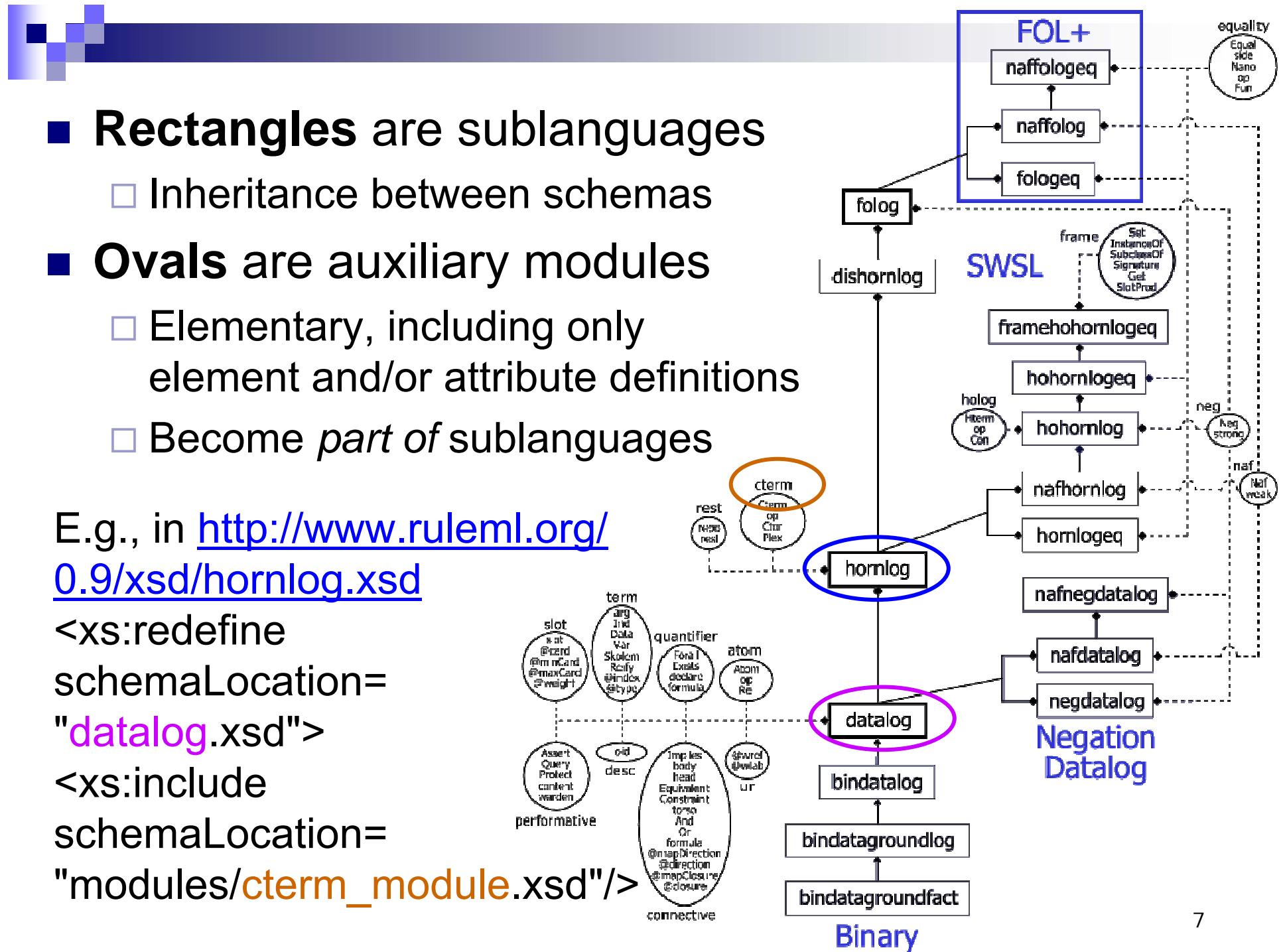
- RuleML: Rule Markup Language
  - RuleML derivation rules (shown here) and production rules defined in XML Schema Definition (XSD)
  - Each XSD of the family corresponds to the expressive class of a specific RuleML sublanguage
- The most recent schema specification of RuleML is always available at <http://www.ruleml.org/spec>
- Current release: RuleML 0.9
- Pre-release: RuleML 0.91

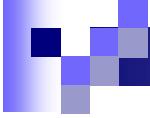


# Schema Modularization

- XSD URIs identify expressive classes
  - Receivers of a rulebase can validate applicability of tools  
(such as **Datalog** vs. **Hornlog** interpreters)
  - Associated with semantic classes  
(such as **function-free** vs. **function-containing** Herbrand models)
- Modularization (Official Model)
  - Aggregation:  
e.g., **Datalog part of Hornlog**
  - Generalization:  
e.g., **Bindatalog is a Datalog**





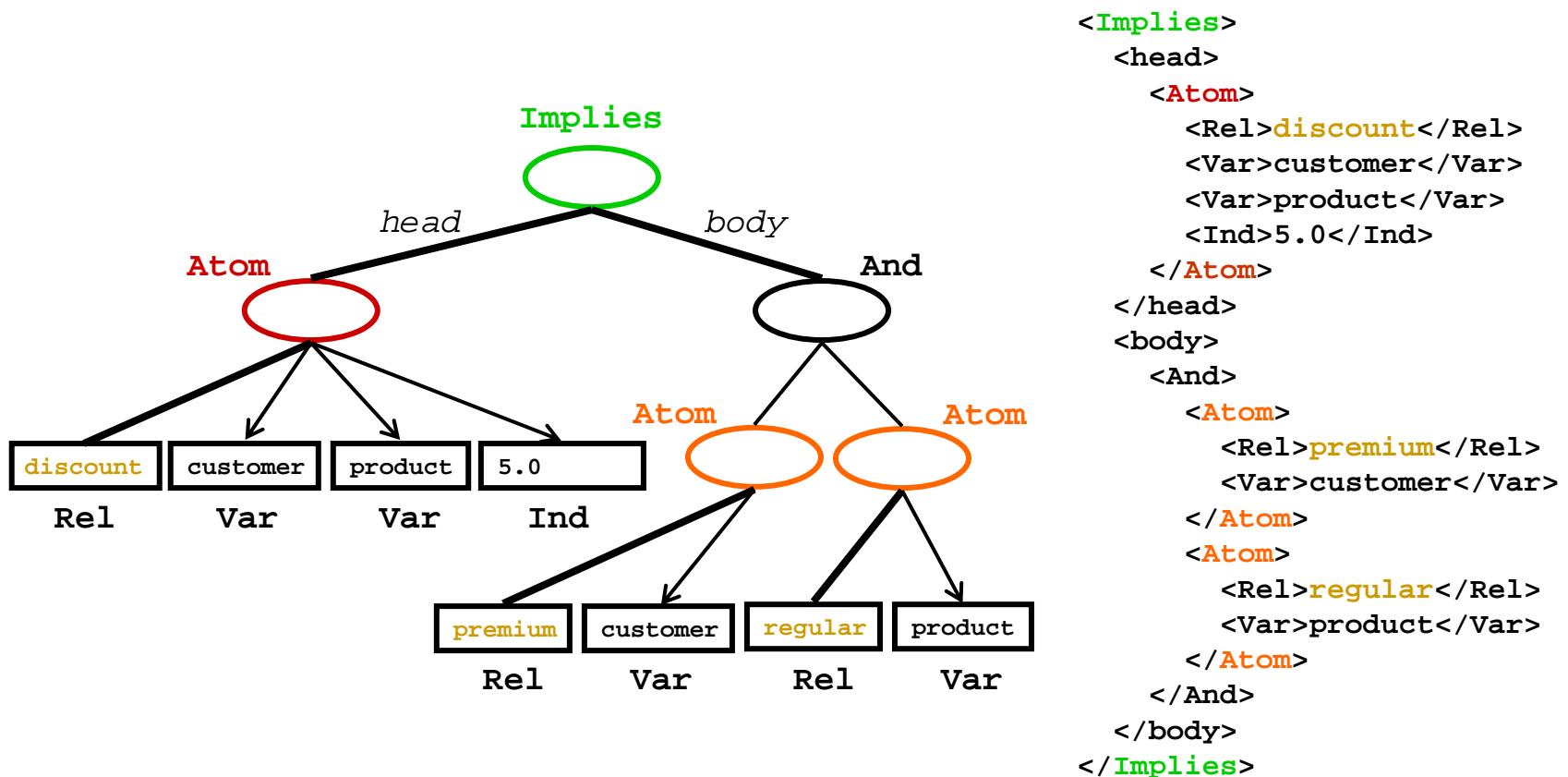


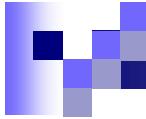
# Bring Datalog to the Semantic Web

- Start with n-ary relations (not binary properties)
- Keep **Variable** typing optional (reuse RDFS' subClassOf taxonomies as sort lattices)
- Allow signature declarations of arities and types
- Employ function-free facts as well as Horn rules (rather than 1<sup>st</sup>: RDF descriptions; 2<sup>nd</sup>: RDF rules)
- Use function-free Herbrand model semantics (querying stays decidable)
- Provide three syntactic levels:
  - User-oriented: Prolog-like, but with “?”-variables
  - Abstract: MOF/UML diagrams
  - XML serialization: Datalog RuleML

# Business Rule: Positional

"The **discount** for a *customer* buying a *product* is 5 percent if the *customer* is **premium** and the *product* is **regular**."



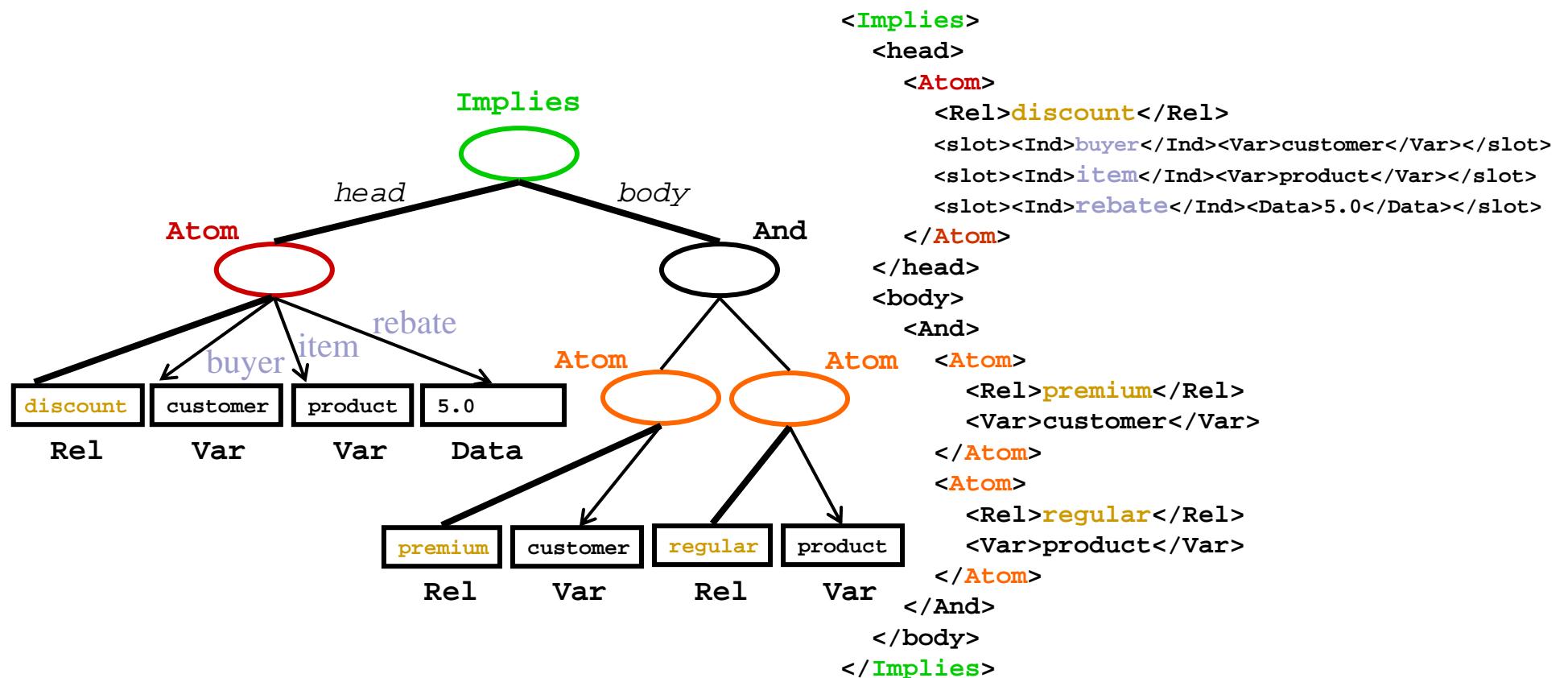


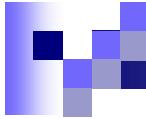
## Extend Datalog for the Semantic Web (I)

- Allow slots as name->filler pairs in **Atoms**  
(cf. F-logic's methods and RDF's properties)
- Extend optional types and signatures for slots
- Add optional object identifiers (**oids**) to atoms
- Separate **Data** literals from **Individual** constants

# Business Rule: Slotted (for OO)

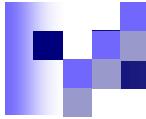
"The **discount** for a *customer* buying a *product* is 5 percent if the *customer* is **premium** and the *product* is **regular**."





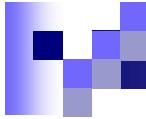
## Extend Datalog for the Semantic Web (II)

- Permit IRI webizing for **Data** (XML Schema Part 2), **Individuals** (RDF's **resources**), **Relations**, **slot** names, types (RDFS' classes), and **oids** (RDF's **about**)
- Introduce **Module** (scope) construct for clauses (cf. RDF's named graphs)
- Add scoped-default (**Naf**), strong (**Neg**), scoped-default-of-strong negation (unscoped: cf. [ERDF](#))
- Integrate with Description Logics
  - Homogeneous (SWRL, Datalog RuleML + OWL-DL)
  - Hybrid (AL-log, Datalog<sup>DL</sup>, DL+log, ...)



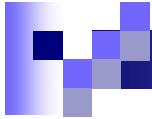
# Bring Horn Logic to the Semantic Web

- Augment Datalog with uninterpreted **Functions** and their **Expressions**; also for extended Datalog
- Augment Datalog's Herbrand model semantics with such **Functions** (querying becomes undecidable)
- Extend Datalog syntaxes
  - XML Schema of Hornlog RuleML inherits and augments XML Schema of Datalog RuleML
- Add **Equality** and **interpreted Functions** (XML serialization: attribute **in="yes"**)
- Reuse XQuery/XPath functions and operators as built-ins



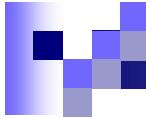
# Specify a First-Order Logic Web Language

- Layer on top of either
  - Disjunctive Datalog: **Or** in the head generalizing Datalog
  - Disjunctive Horn Logic: **Or** in head of near-Horn clauses
- Alternatively, layer on top of either
  - Disjunctive Datalog with restricted strong **Negation**
  - Disjunctive Horn Logic with restricted strong **Neg**
- Permit unrestricted **Or**, **And**, strong **Neg**, and quantifiers **Forall** and **Exists** to obtain FOL
- Use semantics of classical FOL model theory
- Extend Hornlog RuleML syntax to FOL RuleML



# Equality for Functions

- Functional programming (FP) plays increasing Web role: MathML, XSLT, XQuery
- Functional RuleML employs orthogonal notions freely combinable with Relational RuleML
- Also solves a Relational RuleML issue, where the following ‘child-of-parent’ elements are separated:
  - Constructor (**Ctor**) of a complex term (**Cterm**)
  - User-defined function (**Fun**) of a call (**Nano**)
- Proceed to a **logic with equality**



## Function Interpretedness (I)

- Different notions of ‘function’ in LP and FP:
- **LP:** *Uninterpreted functions denote* unspecified values when applied to arguments, not using function definitions
- **FP:** *Interpreted functions compute* specified returned values when applied to arguments, using function definitions
- E.g.: **first-born**: Man × Woman → Human
  - Uninterpreted: **first-born(John, Mary)** denotes first-born
  - Interpreted: using **first-born(John, Mary) = Jory**, so the application returns **Jory**

## Function Interpretedness (II)

- Uninterpreted **<Ctor>** vs. interpreted **<Fun>** functions now distinguished with attribute values:  
**<Fun in="no">** vs. **<Fun in="yes">**
- Function applications with **Cterm** vs. **Nano** then uniformly become **Expressions**
- Two versions of example marked up as follows (where "u" stands for "no" or "yes"):

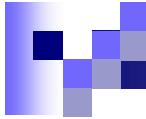
**<Expr>**

**<Fun in="u">first-born</Fun>**

**<Ind>John</Ind>**

**<Ind>Mary</Ind>**

**</Expr>**



# Unconditional Equations

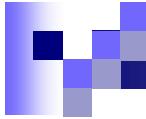
- Modified **<Equal>** element permits both symmetric and oriented equations
- E.g.: **first-born(John, Mary) = Jory** can now be marked up thus:

```
<Equal oriented="yes">
  <lhs>
    <Expr>
      <Fun in="yes">first-born</Fun>
      <Ind>John</Ind>
      <Ind>Mary</Ind>
    </Expr>
  </lhs>
  <rhs>
    <Ind>Jory</Ind>
  </rhs>
</Equal>
```

# Conditional Equations

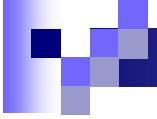
- Use **<Equal>** as the conclusion of an **<Implies>**, whose condition may employ other equations
- E.g.: `?B = birth-year(?P) ⇒ age(?P) = subtract(this-year(),?B)`

```
<Implies>
  <Equal oriented="no">
    <Var>B</Var>
    <Expr>
      <Fun in="yes">birth-year</Fun>
      <Var>P</Var>
    </Expr>
  </Equal>
  <Equal oriented="yes">
    <Expr>
      <Fun in="yes">age</Fun>
      <Var>P</Var>
    </Expr>
    <Expr>
      <Fun in="yes">subtract</Fun>
      <Expr>
        <Fun in="yes">this-year</Fun>
      </Expr>
      <Var>B</Var>
    </Expr>
  </Equal>
</Implies>
```



# Accommodate SWSL-Rules

- HiLog: Higher-order **V**ariables, **C**onstants, and **H**terms (complex terms and atomic formulas at the same time)
- **E**qual: As in Horn Logic with (**unoriented**) **E**quality
- Frames:
  - Value molecules: **A**toms with an **oid**, an optional **R**el class, and zero or more name->filler instance **s**lots
  - Signature molecules: name=>filler class **s**lots, which can have { min : max } cardinality constraints
- Reification: A formula (e.g., a rule) embedded in a **R**eify element is treated (e.g., unified) as a term
- **S**kolems: Unnamed, represent new individual constants (like RDF's blank nodes); otherwise, uniquely named ones



# HiLog Examples: Hterms (I)

- First-order terms:  $f(a, ?x)$

**<Hterm>**

**<op><Con>f</Con></op>**

**<Con>a</Con>**

**<Var>x</Var>**

**</Hterm>**

- Variables over function symbols:  $?X(a, ?Y)$

**<Hterm>**

**<op><Var>x</Var></op>**

**<Con>a</Con>**

**<Var>y</Var>**

**</Hterm>**

## HiLog Examples: Hterms (II)

- Parameterized function symbols:  $f(?X, a)(b, ?X(c))$

```
<Hterm>
```

```
  <op>
```

```
    <Hterm>
```

```
      <op><Con>f</Con></op>
```

```
      <Var>X</Var>
```

```
      <Con>a</Con>
```

```
    </Hterm>
```

```
  </op>
```

```
  <Con>b</Con>
```

```
  <Hterm>
```

```
    <op><Var>X</Var></op>
```

```
    <Con>c</Con>
```

```
  </Hterm>
```

```
</Hterm>
```

# Equality Example

- Equality  $::=$  in rule head:  $f(a,?X)::=g(?Y,b) :- p(?X,?Y).$

```
<Implies>
  <head>
    <Equal>
      <Hterm>
        <op><Con>f</Con></op>
        <Con>a</Con>
        <Var>X</Var>
      </Hterm>
      <Hterm>
        <op><Con>g</Con></op>
        <Var>Y</Var>
        <Con>b</Con>
      </Hterm>
    </Equal>
  </head>
  <body>
    <Hterm>
      <op><Con>p</Con></op>
      <Var>X</Var>
      <Var>Y</Var>
    </Hterm>
  </body>
</Implies>
```

## Frame Example: Value Molecule

- Parameterized-name->filler slot: o[f(a,b) -> 3]

```
<Atom>
  <oid><Con>o</Con></oid>
  <slot>
    <Hterm>
      <op><Con>f</Con></op>
      <Con>a</Con>
      <Con>b</Con>
    </Hterm>
    <Con>3</Con>
  </slot>
</Atom>
```

# Reification Example: Reified Rule

- \$Rule as slot filler: `john[believes -> ${p(?X) implies q(?X)}].`

```
<Hterm>
  <oid>john</oid>
  <slot>
    <Con>believes</Con>
    <Reify>
      <Implies>
        <body>
          <Hterm>
            <op><Con>p</Con></op>
            <Var>X</Var>
          </Hterm>
        </body>
        <head>
          <Hterm>
            <op><Con>q</Con></op>
            <Var>X</Var>
          </Hterm>
        </head>
      </Implies>
    </Reify>
  </slot>
</Hterm>
```

# Skolem Examples (I):

- Named Skolem: `holds(a,_#1)` and `between(1,_#1,5)`.

`<And>`

`<Hterm>`

`<op><Con>holds</Con></op>`

`<Con>a</Con>`

`<Skolem>1</Skolem>`

`</Hterm>`

`<Hterm>`

`<op><Con>between</Con></op>`

`<Con>1</Con>`

`<Skolem>1</Skolem>`

`<Con>5</Con>`

`</Hterm>`

`</And>`

## Skolem Examples (II):

- Unnamed Skolem: `holds(a,_#)` and `between(1,_#,5)`.

`<And>`

`<Hterm>`

`<op><Con>holds</Con></op>`

`<Con>a</Con>`

`<Skolem/>`

`</Hterm>`

`<Hterm>`

`<op><Con>between</Con></op>`

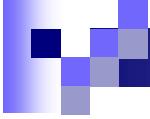
`<Con>1</Con>`

`<Skolem/>`

`<Con>5</Con>`

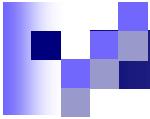
`</Hterm>`

`</And>`



# Proceed to Modal Logics

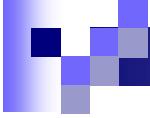
- Modal operators **generically** represented as **Relations** at least one of whose arguments is a proposition represented as an **Atom** with an uninterpreted **Relation** (including another modal operator)
  - *Alethic* **necessary** ( $\Box$ ) and **possible** ( $\Diamond$ )
  - *Deontic* **must** and **may** (e.g., in business rules)
  - Open for *temporal* (e.g., when planning/diagnosing reactive rules), *epistemic* (e.g., in authentication rules), and further modal operators
- Towards a unified framework for multi-modal logic based on Kripke-style possible worlds semantics



# Modal Examples: Alethic Operator

- Necessity:  $\Box \text{prime}(1)$

```
<Atom>  
  <Rel modal="yes">necessary</Rel>  
  <Atom>  
    <Rel in="no">prime</Rel>  
    <Data>1</Data>  
  </Atom>  
</Atom>
```



# Modal Examples: Epistemic Operator

- Knowledge: **knows(Mary,material(moon,rock))**

<Atom>

  <Rel modal="yes">knows</Rel>

  <Ind>Mary</Ind>

  <Atom>

    <Rel in="no">material</Rel>

    <Ind>moon</Ind>

    <Ind>rock</Ind>

  </Atom>

</Atom>



# Modal Examples: Epistemic Reasoning

- Veridicality axiom:  $\text{Knows}_{\text{Agent}} \text{ proposition} \rightarrow \text{proposition}$

$\text{Knows}_{\text{Mary}} \text{material(moon,rock)} \rightarrow \text{material(moon,rock)}$

Serialization in previous slide

→

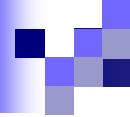
**<Atom>**

```
<Rel in="yes">material</Rel>    <!-- "yes" is default -->
<Ind>moon</Ind>
<Ind>rock</Ind>
</Atom>
```

# Modal Examples: Nested Operators

- Knowledge of Necessity: `knows(Mary, □ prime(1))`

```
<Atom>  
  <Rel modal="yes">knows</Rel>  
  <Ind>Mary</Ind>  
  <Atom>  
    <Rel modal="yes" in="no">necessary</Rel>  
    <Atom>  
      <Rel in="no">prime</Rel>  
      <Data>1</Data>  
    </Atom>  
  </Atom>  
</Atom>
```



# Protect Knowledge Bases by Integrity Constraints

- A knowledge base KB is a formula in any of our logic languages
- An integrity constraint IC is also a formula in any of our logic languages, which may be chosen independently from KB
- KB obeys IC  
    iff

**KB entails IC**

(Reiter [1984](#), [1987](#))

- Entailment notion of [1987](#) uses epistemic modal operator
- Serialization: <Protect> IC KB </Protect>

# Integrity Constraint Example: Rule with $\exists$ -Head

- Adapted from (Reiter 1987):

$$IC = \{(\forall x) \text{emp}(x) \Rightarrow (\exists y) \text{ssn}(x, y)\}$$

$KB_1 = \{\text{emp}(\text{Mary})\}$   $KB_1$  violates IC

$KB_2 = \{\text{emp}(\text{Mary}), \text{ssn}(\text{Mary}, 1223)\}$   $KB_2$  obeys IC

<Protect> IC  $KB_i$  </Protect>

IC:

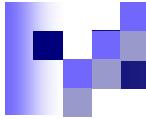
```
<Forall>
  <Var>x</Var>
  <Implies>
    <Atom>
      <Rel>emp</Rel>
      <Var>x</Var>
    </Atom>
    <Exists>
      <Var>y</Var>
      <Atom>
        <Rel>ssn</Rel>
        <Var>x</Var>
        <Var>y</Var>
      </Atom>
    </Exists>
  </Implies>
</Forall>
```

$KB_1$ :

```
<Atom>
  <Rel>emp</Rel>
  <Ind>Mary</Ind>
</Atom>
```

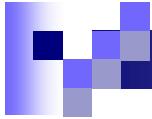
$KB_2$ :

```
<And>
  <Atom>
    <Rel>emp</Rel>
    <Ind>Mary</Ind>
  </Atom>
  <Atom>
    <Rel>ssn</Rel>
    <Ind>Mary</Ind>
    <Data>1223</Data>
  </Atom>
</And>
```



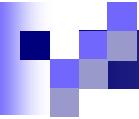
# Approach Production and Reaction Rules

- Share Condition (C) part with earlier languages as proposed for the [RIF Condition Language](#)
- Develop Action (A) part of Production Rules via a taxonomy of actions on KBs (Assert, Retract, ...), on local or remote hosts, or on the surroundings
- Develop Event (E) part of Reaction Rules via a corresponding taxonomy
- Create CA and ECA families bottom-up and map to relevant languages for Semantic Web Services
- Serialization: <**React**> E C A </**React**>



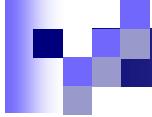
# RDF Rules

- RDF-like Rules: Important RuleML sublanguage
  - Datalog: Relational databases augmented by views
  - RDF Properties: Slots permit non-positional, keyed arguments
  - RDF URIs: Anchors provide **O**bject **i**dentity via webzing through URIs
    - **oids**: Can be **I**ndividuals, **V**ariables, etc.
    - **uris**: Now used for both RDF's **about** and **resource**
  - RDF Blank Nodes: F-logic/Flora-2 Skolem-constant approach
    - E.g., Skolem generator '\_' becomes <**skolem**/>



```
<Implies>
  <body>
    <And>
      <Atom>
        <oid><Var>x</Var></oid>
        <Rel>product</Rel>
        <slot><Ind uri=":price"/><Var>y</Var></slot>
        <slot><Ind uri=":weight"/><Var>z</Var></slot>
      </Atom>
      <Atom>
        <Rel uri="swrlb:greaterThan"/><Var>y</Var><Data>200</Data>
      </Atom>
      <Atom>
        <Rel uri="swrlb:lessThan"/><Var>z</Var><Data>50</Data>
      </Atom>
    </And>
  </body>
  <head>
    <Atom>
      <oid><Var>x</Var></oid>
      <Rel>product</Rel>
      <slot><Ind uri=":shipping"/><Data>0</Data></slot>
    </Atom>
  </head>
</Implies>
```

**“For a product whose price is greater than 200 and whose weight is less than 50, no shipping is billed.”**



# Bidirectional Interpreters in Java

- Two varieties of reasoning engines
  - **Top-Down**: backward chaining
  - **Bottom-Up**: forward chaining
- jDREW: *Java Deductive Reasoning Engine for the Web* includes both TD and BU  
<http://www.jdrew.org>
- OO jDREW: *Object-Oriented* extension to jDREW  
<http://www.jdrew.org/oojdrew>
- Java Web Start online demo available at  
<http://www.jdrew.org/oojdrew/demo.html>

# OO jDREW Slots

- Normalized atoms and complex terms
  - **oids** (object identifier)
  - **Positional** parameters (in their original order)
  - **Positional** rest terms
  - **Slotted** parameters (in the order encountered)
  - **Slotted** rest terms
- Efficient unification algorithm
  - Linear  $O(m+n)$ : instead of  $O(m^*n)$ 
    - No need for positional order
    - Slots internally sorted
  - Steps:
    - Scan two lists of parameters
      - Matching up roles and positions for positional parameters
      - Unifying those parameters
    - Add unmatched roles to list of rest terms
    - Generate dynamically a Plex (RuleML's closest equivalent to a list) for a collection of rest terms

OO jDREW Top-Down Engine

Type Definition Knowledge Base Query

Query:  
discount(?person, ?thing, ?amount).

Solution:  
\$top0:-discount(PeterMiller, Honda, percent5).  
?- discount(PeterMiller, Honda, percent5):-premium(PeterMiller),regular(Honda).  
    |- premium(PeterMiller).  
    |- regular(Honda).

Variable Bindings:

Variable	Binding
?person	PeterMiller
?thing	Honda
?amount	percent5

Issue Query Next Solution

Show Debug Console

Java Application Window

**discount(?customer,?product,percent5)  
:- premium(?customer), regular(?product).**

**premium(PeterMiller).  
regular(Honda).**

positional

OO jDREW Top-Down Engine

Type Definition Knowledge Base Query

Query:

```
discount(rebate->?amount;prod->?thing;cust->?person).
```

Solution:

```
$top() :- discount(cust->PeterMiller; prod->Honda; rebate->percent5).  
?- discount(cust->PeterMiller; prod->Honda; rebate->percent5) :- premium(cust->PeterMiller).  
    regular(prod->Honda).
```

Issue Query Next Solution

Variable Bindings:

Variable	Binding
?person	cust->PeterMiller
?thing	prod->Honda
?amount	rebate->percent5

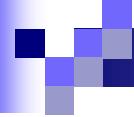
Show Debug Console

Java Application Window



**discount(cust->?customer;prod->?product;rebate->percent5)  
:- premium(cust->?customer), regular(prod->?product).**

**premium(cust->PeterMiller).  
regular(prod->Honda).**



# OO jDREW Types

- Order-sorted type system
  - RDF Schema: lightweight taxonomies of the Semantic Web
  - To specify a partial order for a set of classes in RDFS
- Advantages
  - Having the appropriate types specified for the parameters
  - To restrict the search space
  - Faster and more robust system than when reducing types to unary predicate calls in the body
- Limitations
  - Only modeling the taxonomic relationships between classes
  - Not modeling properties with domain and range restrictions

OO jDREW Top-Down Engine

Type Definition Knowledge Base Query

Query:

```
base_price(customer->[sex->male; name->"John Doe"; age->28]; vehicle->vehicle:ToyotaCorolla; price->?money:Integer).
```

Solution:

```
$top() :- base_price(customer->[sex->male; name->"John Doe"; age->28];
  base_price(customer->[sex->male; name->"John Doe"; age->28!];
```

Variable Bindings:

Variable	Binding
?money : Integer	price->650 : Integer

**base\_price(customer->[sex->male; !];  
vehicle->:Car;  
price->650:Integer).**

**base\_price(customer->[sex->male; !];  
vehicle->:Van;  
price->725:Integer).**

```

graph TD
    Thing --> Vehicle
    Vehicle --> PassengerVehicle
    Vehicle --> Van
    PassengerVehicle --> Car
    PassengerVehicle --> MiniVan
    Car --> Sedan1
    Car --> Sedan2
    ToyotaCorolla --> Sedan1
    Nothing --> Sedan1
    Nothing --> Sedan2
    
```

The diagram illustrates an object-oriented type hierarchy. At the top is the abstract type **Thing**, which has a single inheritance arrow pointing down to **Vehicle**. From **Vehicle**, two arrows point down to **PassengerVehicle** and **Van**. From **PassengerVehicle**, two arrows point down to **Car** and **MiniVan**. Finally, from **Car**, two arrows point down to **Sedan** (labeled "Sedan") and **Sedan** (labeled "ToyotaCorolla"). Additionally, there are two dashed arrows originating from a box labeled **Nothing**, one pointing to each of the two **Sedan** nodes.

Java Application Window

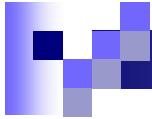
# OO jDREW OIDs

- **oid**: Object Identifier
- Currently: symbolic names
  - In **<Atom>** & **<Implies>**
- Planned: **uri** attribute
- E.g., give name to fact  
keep(Mary, ?object).

```
<Atom>
<oid><Ind>mary-12</Ind></oid>
<Rel>keep</Rel>
<Ind>Mary</Ind>
<Var>object</Var>
</Atom>

<Atom>
<oid><Ind uri="http://mkb.ca"/></oid>
<Rel>keep</Rel>
<Ind>Mary</Ind>
<Var>object</Var>
</Atom>

<Atom>
<oid><Var>object</Var></oid>
<Rel>keep</Rel>
<Ind>Mary</Ind>
<Var>object</Var>
</Atom>
```



## Conclusions

- RuleML is modular family, whose root allows to access the language as a whole and whose members allow customized subsets
- New members joining, e.g. Fuzzy RuleML
- Concrete & abstract syntax of RuleML
  - Specified by modular XSD (shown here) & MOF
- Formal semantics of OO Hornlog RuleML
  - Implemented by OO jDREW BU & TD
- Interoperability/Interchange of/with RuleML
  - Realized by translators, primarily via XSLT