

# Process-Algebra Based Query Workflows

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# Motivating Example: Travel Planning

- *Query:* What is the cheapest connection from Freiburg to Amsterdam?
- *Solution:*
  - ▶ Consider nearest airports
  - ▶ Query train connections from start to airports
  - ▶ Query flights from airports to destination
  - ▶ ...

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→ Design a *Query Workflow* to find connections *automatically*

# The MARS Framework

- Modular Active Rules for the Semantic Web
- Provides an open framework for active rules and for processes
- Supports *heterogeneous event, query, and action* languages:
  - ▶ XQuery, SPARQL
  - ▶ WSQL for querying Web Services
  - ▶ DWQL for querying Deep Web sources
- *Data model*: Relational dataflow

Set of tuples of variable bindings  $R$ , i.e., every tuple is of the form  $t = \{v_1/x_1, \dots, v_n/x_n\}$  with  $v_1, \dots, v_n$  variables and  $x_1, \dots, x_n$  elements of the underlying domain

## Data Model: Example

Consider the query:

$$\text{Query}((dt, at, pr) \leftarrow \text{getTrainConnection}(start, dest, date))$$

which returns the departure time ( $dt$ ), arrival time ( $at$ ), and price ( $pr$ ) for a train connection from  $start$  to  $dest$  on  $date$ :

```
getTrainConnection(  
  {{start/"FR", dest/"BS", date/"10.06.2009"},  
   {start/"FR", dest/"F", date/"10.06.2009"}}) =  
  {{start/"FR", dest/"BS", date/"10.06.2009",  
   dt/"10:02", at/"10:47", pr/"21.80"},  
   ... ,  
  {start/"FR", dest/"F", date/"10.06.2009",  
   dt/"09:57", at/"12:43", pr/"61.00"}}
```

# Calculus of Communicating Systems (CCS)

- Process Algebra
- Express processes over atomic constituents by:
  - ▶ Sequence
  - ▶ Alternative
  - ▶ Concurrent
  - ▶ Recursion
- Process state is encoded implicitly in the process itself

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→ RelCCS: Extension of CCS with relational dataflow

# Atomic Constituents of RelCCS

## Formal Semantics of Process Execution

$$\llbracket \cdot \rrbracket : \mathcal{P} \times 2^{\text{Tuples}(\text{Var})} \rightarrow 2^{\text{Tuples}(\text{Var})}$$

- *Actions*:  $\llbracket \text{Action}(a)[R] \rrbracket := R$ , plus external side effects of  $a$
- *Queries*:  $\llbracket \text{Query}(q)[R] \rrbracket := \bigcup_{t \in R} \llbracket \text{Query}(q)[t] \rrbracket = R \bowtie q_0$
- *Tests*:  $\llbracket \text{Test}(c)[R] \rrbracket = \sigma[c](R)$ , like SQL's  
SELECT \* FROM R WHERE  $c$
- *Events*: Similar to queries. On occurrence of an event, matching tuples proceed through the process.

# Selected Operators of RelCCS

- *Sequence*:  $\llbracket \text{Seq}(P, Q)[R] \rrbracket := \llbracket Q[\llbracket P[R] \rrbracket] \rrbracket$
- *Union*: Each branch is started with  $R$ .  
 $\llbracket \text{Union}(P_1, \dots, P_n)[R] \rrbracket = \llbracket P_1[R] \rrbracket \cup \dots \cup \llbracket P_n[R] \rrbracket$
- *Concurrent*: Each branch is started with  $R$ .  
 $\llbracket \text{Concurrent}(P_1, \dots, P_n)[R] \rrbracket := \llbracket P_1[R] \rrbracket \bowtie \dots \bowtie \llbracket P_n[R] \rrbracket$

# Application Scenario: Travel Planning

```
# process input: (start, dest, date)
Sequence(
  Query(rd ← distance(start, dest)),
  Union(
    Sequence(Test(rd < 800), # consider to go by train
      Query((dt, at, pr) ← getTrainConnection(start, dest, date))),
    Sequence(Test(rd ≥ 400), # consider also to use flights
      Query(ap ← getAirports()),
      Query(dist ← distance(start, ap)),
      TopK(10,100,null,dist,xsd:decimal,asc,false),
      Query((dt, at, pr) ← getTrainConnection(start, ap, date)),
      UseDefinition(exploreTC[ap])))
```

**Definition(exploreTC[local: ap, ...]) := ...**

# postcondition: reached, either by train or train+flight<sup>+</sup>, or  
train+flight<sup>+</sup>+train

# Related Work

- Petri Nets
  - ▶ Graphical formalism with concise formal semantics
  - ▶ Support for formal analysis and verification techniques
- Yet Another Workflow Language (YAWL)
  - ▶ Based on exhaustive analysis of workflow patterns
  - ▶ Roots in Petri Nets
- WS-BPEL
  - ▶ Block-oriented process language, similar to procedural programming languages
  - ▶ Dataflow described by variables, that can be made set-oriented by referencing external database tables

# Conclusion

- RelCCS: Extension of CCS with relational dataflow
  - ▶ Both the primitives for control structures and data flow are on the same level of the language
  - ▶ Embedding in MARS framework provides tight integration, e. g. RelCCS fragments can be used as action parts in MARS' active rules
  - ▶ Comfortable embedding of algorithmic Web Services (Configurable Graph Datatype)
- *Future Work*: Automatic derivation of processes and patterns from ontologies

## Prototype Implementation

RelCCS is implemented in a prototype which can be found with sample processes and further documentation at

<http://www.semwebtech.org/mars/frontend/> → run CCS Process.

# ??? Questions ???

## Take-Home Message

Often it is easier to design the process *how* to solve a problem than stating a single query