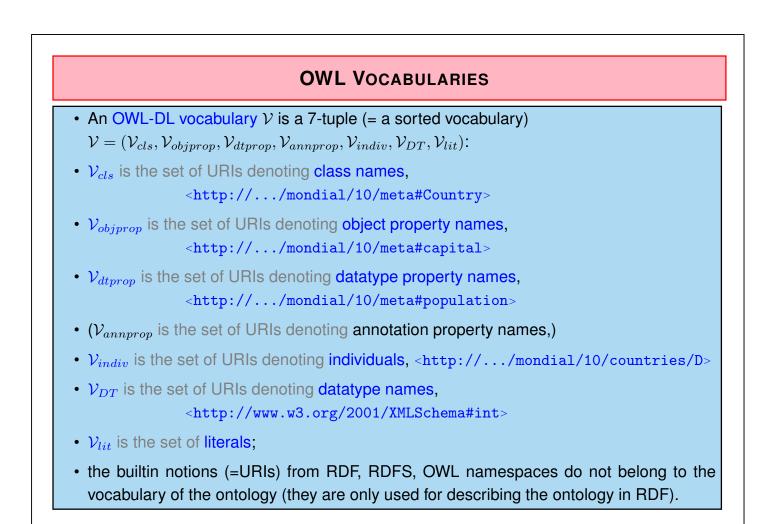
### 7.2 OWL

- the OWL versions use certain DL semantics:
- Base:  $ALC_{R+}$ : (i.e., with transitive roles). This logic is called S (reminiscent to its similarity to the modal logic S).
- roles can be ordered hierarchically (rdfs:subPropertyOf;  $\mathcal{H}$ ).
- OWL Lite: SHIF(D), Reasoning in EXPTIME.
- OWL DL: SHOIN(D), decidable. Pellet (2007) implements SHOIQ(D). Decidability is in NEXPTIME (combined complexity wrt. TBox+ABox), but the actual complexity of a given task is constrained by the maximal used cardinality and use of nominals and inverses and behaves like the simpler classes. (Ian Horrocks and Ulrike Sattler: A Tableau Decision Procedure for SHOIQ(D); In IJCAI, 2005, pp. 448-453; available via http://dblp.uni-trier.de)
- OWL 2.0 towards SROIQ(D) and more datatypes ...

# **OWL NOTIONS; OWL-DL VS. RDF/RDFS; MODEL VS. GRAPH**

- OWL is defined based on (Description Logics) model theory,
- OWL ontologies can be represented by RDF graphs,
- Only certain RDF graphs are allowed OWL-DL ontologies: those, where class names, property names, individuals etc. occur in a well-organized way.
- · Reasoning works on the (Description Logic) model, the RDF graph is only a means to represent it.

(recall: RDF/RDFS "reasoning" works on the graph level)



312

# **OWL INTERPRETATIONS**

Since DL is a subset of FOL, the interpretation of an OWL-DL vocabulary can be given as a FOL interpretation

 $\mathcal{I} = (I_{indiv} \cup I_{cls} \cup I_{objprop} \cup I_{dtprop} \cup I_{annprop} \cup I_{DT} , \mathcal{U}_{obj} \cup \mathcal{U}_{DT})$ 

where I interprets the vocabulary as

- *I*<sub>indiv</sub> constant symbols (individuals),
- $I_{cls}$ ,  $I_{DT}$  unary predicates (classes and datatypes),
- *I*<sub>objprop</sub>, *I*<sub>dtprop</sub>, *I*<sub>annprop</sub> binary predicates (properties),

and the universe  $\ensuremath{\mathcal{U}}$  is partitioned into

- an object domain  $\mathcal{U}_{obj}$
- and a data domain  $\mathcal{U}_{DT}$  (of all values of datatypes).

### **OWL INTERPRETATIONS**

The interpretation I is as follows:

$I_{indiv}$ :	each individual $a \in \mathcal{V}_{indiv}$ to an object $I(a) \in \mathcal{U}_{obj}$ ,
	(e.g., $I(<\texttt{http:///mondial/10/countries/D}>) = germany$ )
$I_{cls}$ :	each class $C \in \mathcal{V}_{cls}$ to a set $\ I(C) \subseteq \mathcal{U}_{obj},$
	$(\texttt{e.g.}, \textit{germany} \in I(\texttt{}))$
$I_{DT}$ :	each datatype $D \in \mathcal{V}_{DT}$ to a set $I(D) \subseteq \mathcal{U}_{DT}$ ,
	(e.g., $I(<\texttt{http://www.w3.org/2001/XMLSchema#int}) = \{\dots, -2, -1, 0, 1, 2, \dots\}$ )
$I_{objprop}$ :	each object property $p\in\mathcal{V}_{objprop}$ to a binary relation $\ I(p)\subseteq\mathcal{U}_{obj} imes\mathcal{U}_{obj}$ ,
	$(\texttt{e.g.}, (\textit{germany}, \textit{berlin}) \in I(\texttt{}))$
$I_{dtprop}$ :	each datatype property $p\in\mathcal{V}_{dtprop}$ to a binary relation $\ I(p)\subseteq\mathcal{U}_{obj} imes\mathcal{U}_D$ ,
	$(\texttt{e.g.}, (\textit{germany}, 83536115) \in I(\texttt{}))$
$I_{annprop}$ :	each annotation property $p\in\mathcal{V}_{annprop}$ to a binary relation $\ I(p)\subseteq\mathcal{U} imes\mathcal{U}.$

314

### OWL Class Definitions and Axioms (Overview)

- owl:Class
- The properties of an owl:Class (including owl:Restriction) node describe the properties of that class.

An owl:Class is required to satisfy the conjunction of all constraints (implicit: intersection) stated about it.

These characterizations are roughly the same as discussed for DL class definitions:

- Constructors: owl:unionOf, owl:intersectionOf, owl:complementOf (ALC)
- Enumeration Constructor: owl:oneOf (enumeration of elements;  $\mathcal{O}$ )
- Axioms rdfs:subClassOf, owl:equivalentClass,
- Axiom owl:disjointWith (also expressible in ALC: *C* disjoint with *D* is equivalent to  $C \sqsubseteq \neg D$ )

# OWL NOTIONS (CONT'D)

OWL Restriction Classes (Overview)

- owl:Restriction is a subclass of owl:Class, allowing for specification of a constraint on one property.
- one property is restricted by an owl:onProperty specifier and a constraint on this property:
  - $(\mathcal{N}, \mathcal{Q}, \mathcal{F})$  owl:cardinality, owl:minCardinality or owl:maxCardinality,
  - owl:allValuesFrom ( $\forall R.C$ ), owl:someValuesFrom ( $\exists R.C$ ),
  - owl:hasValue ( $\mathcal{O}$ ),
  - including datatype restrictions for the range (D)
- by defining intersections of owl:Restrictions, classes having multiple such constraints can be specified.

316

# OWL NOTIONS (CONT'D)

**OWL** Property Axioms (Overview)

- Distinction between owl:ObjectProperty and owl:DatatypeProperty
- from RDFS: rdfs:domain/rdfs:range assertions, rdfs:subPropertyOf
- Axiom owl:equivalentProperty
- Axioms: subclasses of rdf:Property: owl:TransitiveProperty, owl:SymmetricProperty, owl:FunctionalProperty, owl:InverseFunctionalProperty (see Slide 332)

OWL Individual Axioms (Overview)

- · Individuals are modeled by unary classes
- owl:sameAs, owl:differentFrom, owl:AllDifferent(o<sub>1</sub>,...,o<sub>n</sub>).

# FIRST-ORDER LOGIC EQUIVALENTS

$OWL: x \in C$	DL Syntax	FOL
С	C	C(x)
intersectionOf $(C_1, C_2)$	$C_1 \sqcap \ldots \sqcap C_n$	$C_1(x) \wedge \ldots \wedge C_n(x)$
$unionOf(C_1,C_2)$	$C_1 \sqcup \ldots \sqcup C_n$	$C_1(x) \lor \ldots \lor C_n(x)$
$complementOf(C_1)$	$\neg C_1$	$\neg C_1(x)$
$oneOf(x_1,\ldots,x_n)$	$\{x_1\}\sqcup\ldots\sqcup\{x_n\}$	$x = x_1 \lor \ldots \lor x = x_n$
$OWL: x \in C$ , Restriction on $P$	DL Syntax	FOL
someValuesFrom(C')	$\exists P.C'$	$\exists y: P(x,y) \wedge C'(y)$
allValuesFrom $(C')$	$\forall P.C'$	$\forall y: P(x,y) \to C'(y)$
hasValue(y)	$\exists P.\{y\}$	P(x,y)
maxCardinality(n)	$\leq n.P$	$\exists^{\leq n}y:P(x,y)$
minCardinality(n)	$\geq n.P$	$\exists^{\geq n}y:P(x,y)$
cardinality(n)	n.P	$\exists^{=n}y:P(x,y)$

First-	ORDER LOGIC E	QUIVALENTS (CO
OWL Class Axioms for $C$	DL Syntax	FOL
rdfs:subClassOf( $C_1$ )	$C \sqsubseteq C_1$	$\forall x : C(x) \to C_1(x)$
equivalentClass $(C_1)$	$C \equiv C_1$	$\forall x : C(x) \leftrightarrow C_1(x)$
disjointWith $(C_1)$	$C \sqsubseteq \neg C_1$	$\forall x: C(x) \to \neg C_1(x)$
OWL Individual Axioms	DL Syntax	FOL
$x_1$ sameAs $x_2$	$\{x_1\} \equiv \{x_2\}$	$x_1 = x_2$
$x_1$ differentFrom $x_2$	$\{x_1\} \sqsubseteq \neg \{x_2\}$	$x_1 \neq x_2$
AllDifferent $(x_1, \ldots, x_n)$	$\bigwedge_{i\neq j} \{x_i\} \sqsubseteq \neg \{x_j\}$	$\bigwedge_{i \neq j} x_i \neq x_j$

# FIRST-ORDER LOGIC EQUIVALENTS (CONT'D)

OWL Properties	DL Syntax	FOL
Р	P	P(x,y)
OWL Property Axioms for P	DL Syntax	FOL
rdfs:range $(C)$	$\top \sqsubseteq \forall P.C$	$\forall x, y : P(x, y) \to C(y)$
rdfs:domain(C)	$C \sqsupseteq \exists P.\top$	$\forall x, y: P(x, y) \to C(x)$
$subPropertyOf(P_2)$	$P \sqsubseteq P_2$	$\forall x, y : P(x, y) \to P_2(x, y)$
equivalentProperty $(P_2)$	$P \equiv P_2$	$\forall x, y : P(x, y) \leftrightarrow P_2(x, y)$
$inverseOf(P_2)$	$P \equiv P_2^-$	$\forall x, y : P(x, y) \leftrightarrow P_2(y, x)$
TransitiveProperty	$P^+ \equiv P$	$\forall x, y, z : ((P(x, y) \land P(y, z)) \to P(x, z))$
		$\forall x, z : ((\exists y : P(x, y) \land P(y, z)) \to P(x, z))$
FunctionalProperty	$\top \sqsubseteq \leq 1P.\top$	$\forall x, y_1, y_2 : P(x, y_1) \land P(x, y_2) \to y_1 = y_2$
InverseFunctionalProperty	$\top \sqsubseteq {\leq} 1P^{-}.\top$	$\forall x, y_1, y_2 : P(y_1, x) \land P(y_2, x) \rightarrow y_1 = y_2$

### 320

### SYNTACTICAL REPRESENTATION

- OWL specifications can be represented by graphs: OWL constructs have a straightforward representation as triples in RDF/XML and N3.
- there are several logic-based representations (e.g. *Manchester OWL Syntax*); TERP (which can be used with pellet) is a combination of Turtle and Manchester syntax.
- OWL in RDF/XML format: usage of class, property, and individual names:
  - as @rdf:about when used as identifier of a subject (owl:Class, rdf:Property and their subclasses),
  - as @rdf:resource as the object of a property.
- some constructs need auxiliary structures (collections): owl:unionOf, owl:intersectionOf, and owl:oneOf are based on Collections
  - representation in RDF/XML by rdf:parseType="Collection".
  - representation in N3 by  $(x_1 x_2 \dots x_n)$
  - as RDF lists: rdf:List, rdf:first, rdf:rest

# REQUIREMENT

every entity in an OWL ontology must be explicitly typed (i.e., as a class, an object property, a datatype property, ..., or an instance of some class).
 (for reasons of space this is not always done in the examples; in general, it may lead to incomplete results)

322

# QUERYING OWL DATA

- queries are atomic and conjunctive DL queries against the underlying OWL-DL model.
- this model can still be seen as a graph:
  - many of the edges are those known from the basic RDF graph
  - some edges (and collections) are only there for encoding OWL stuff (describing owl:unionOf, owl:propertyChain etc.) – these should not be queried
- SPARQL-DL is a subset of SPARQL: not every SPARQL query pattern is allowed for use on an OWL ontology (but the reasonable ones are, so in practice this is not a problem.)
- the query language SPARQL-DL allows exactly such well-sorted patterns using the notions of OWL.

### SOME TBOX-ONLY REASONING EXAMPLES ON SETS

Example: A Simple Paradox

@prefix : <foo://bla/>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
:Paradox owl:complementOf :Paradox.

[Filename: RDF/paradox.n3]

· without reasoner:

jena -t -ol rdf/xml -if paradox.n3

Outputs the same RDF facts in RDF/XML without checking consistency.

• with reasoner:

jena -e -pellet -if paradox.n3

reads the RDF file, creates a model (and checks consistency) and in this case reports that it is not consistent:

"There is an anonymous individual which is forced to belong to class foo://bla/Paradox and its complement"

• Note: the reasoner invents an anonymous individual for checking consistency. The empty interpretation (with empty domain!) would be a model of  $P \equiv \neq P$ .

324

# Union as $A \sqcup B \equiv \neg((\neg A) \sqcap (\neg B))$ (De Morgan's Rule)

<pre>@prefix : <foo: bla=""></foo:>.</pre>	
<pre>@prefix rdf: <http: 02="" 1999="" 22<="" pre="" www.w3.org=""></http:></pre>	2-rdf-syntax-ns#>.
<pre>@prefix owl: <http: 07="" 2002="" ow<="" pre="" www.w3.org=""></http:></pre>	1#>.
:A rdf:type owl:Class. :B rdf:type	owl:Class.
:Union1 owl:equivalentClass [ owl:unionOf	(:A :B)].
:CompA owl:complementOf :A. :CompB owl:c	complementOf :B.
:IntersectComps owl:equivalentClass [ owl:	<pre>intersectionOf (:CompA :CompB).]</pre>
:Union2 owl:complementOf :IntersectComps.	
:x rdf:type :A. :x rdf:type	:B.
:y rdf:type :CompA. # a negative assertion	y not in A would be better -> OWL 2
:y rdf:type :CompB.	[Filename: RDF/union.n3]
prefix owl: <http: 07="" 2002="" owl<="" td="" www.w3.org=""><td>#&gt;</td></http:>	#>
prefix rdf: <http: 02="" 1999="" 22-<="" td="" www.w3.org=""><td>rdf-syntax-ns#&gt;</td></http:>	rdf-syntax-ns#>
prefix : <foo: bla=""></foo:>	
select ?X ?C ?D	
<pre>from <file:union.n3></file:union.n3></pre>	[Filename: RDF/union.sparql]
where {{?X rdf:type ?C} UNION {:Union1 owl	:equivalentClass ?D}}

### EXAMPLE: UNION AND SUBCLASS

```
@prefix : <foo://bla/>.
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
:Male a owl:Class. ## if these lines are missing,
:Female a owl:Class. ## the reasoner complains
:Person owl:equivalentClass [ owl:unionOf (:Male :Female) ].
:EqToPerson owl:equivalentClass [ owl:unionOf (:Female :Male) ].
:unknownPerson a [ owl:unionOf (:Female :Male) ]. [Filename: RDF/union-subclass.n3]
```

• print class tree (with jena -e -pellet -if union-subclass.n3):

```
owl:Thing
bla:Person = bla:EqToPerson - (bla:unknownPerson)
bla:Female
bla:Male
```

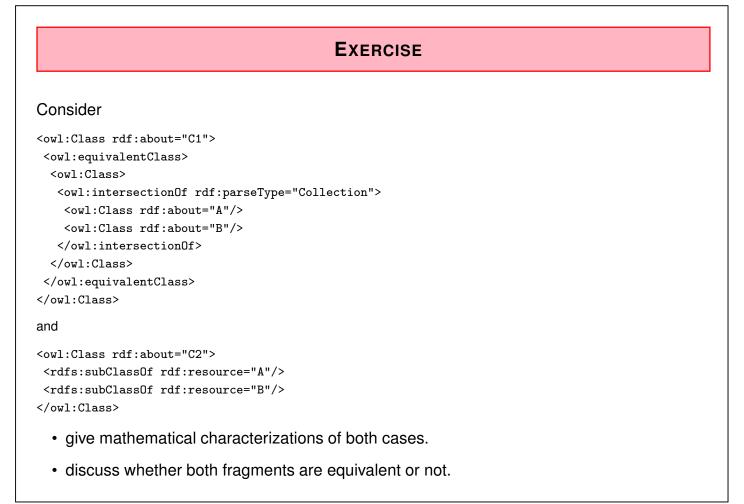
- Male and Female are derived to be subclasses of Person.
- Person and EqToPerson are equivalent classes.
- unknownPerson is a member of Person and EqToPerson.

326

Example (Cont'd)	
prefix rdf: <http: 02="" 1999="" 22-rd:<="" td="" www.w3.org=""><td>f-syntax-ns#&gt;</td></http:>	f-syntax-ns#>
prefix rdfs: <http: 01="" 2000="" rdf-s<="" td="" www.w3.org=""><td>schema#&gt;</td></http:>	schema#>
<pre>prefix owl: <http: 07="" 2002="" owl#="" www.w3.org=""></http:></pre>	
<pre>prefix : <foo: bla=""></foo:></pre>	
select ?SC ?C ?T ?CC ?CD	
<pre>from <file:union-subclass.n3></file:union-subclass.n3></pre>	
where {{?SC rdfs:subClassOf ?C} UNION	
{:unknownPerson rdf:type ?T} UNION	
{?CC owl:equivalentClass ?CD}}	[Filename: RDF/union-subclass.sparql]

• Note: OWLizations of DL class expressions are always handled as blank nodes, and used with "owl:equivalentClass", "rdf:subClassOf", "rdfs:domain", "rdfs:range" or "a".





### DISCUSSION

- Two classes are *equivalent* (wrt. the knowledge base) if they have the same interpretation in every *model* of the KB.
- $C_1$  is characterized to be the intersection of classes A and B.
- for  $C_2$ , it is asserted that  $C_2$  is a subset of A and that it is a subset of B.
- Thus there can be some c that is in A, B,  $C_1$ , but not in  $C_2$ .
- Thus,  $C_1$  and  $C_2$  are not equivalent.
- $C_1$  is a definition, the statements about  $C_2$  are just two constraints ( $C_2$  might be empty).

330

### **DISCUSSION: FORMAL NOTATION**

The DL equivalent to the knowledge base (TBox) is

$$\mathcal{T} = \{ C_1 \equiv (A \sqcap B) , \quad C_2 \sqsubseteq A , \quad C_2 \sqsubseteq B \}$$

The First-Order Logic equivalent is

$$\mathcal{KB} = \{ \forall x : A(x) \land B(x) \leftrightarrow C_1(x) , \quad \forall x : C_2(x) \to A(x) \land B(x) \}$$

Thus,  $\mathcal{KB} \models \forall x : C_2(x) \to A(x) \land B(x)$ .

Or, in DL:  $\mathcal{T} \models C_2 \sqsubseteq C_1$ .

On the other hand,  $\mathcal{M} = (\mathcal{D}, \mathcal{I})$  with  $\mathcal{D} = \{c\}$  and

 $\mathcal{I}(A) = \{c\}, \ \mathcal{I}(B) = \{c\}, \ \mathcal{I}(C_1) = \{c\}, \ \mathcal{I}(C_2) = \emptyset$ 

is a model of  $\mathcal{KB}$  (wrt. first-order logic) and  $\mathcal{T}$  (wrt. DL) that shows that  $C_1$  and  $C_2$  are not equivalent.

# SUBCLASSES OF PROPERTIES Triple syntax: some property rdf:type a specific type of property According to their ranges • owl:ObjectProperty – subclass of rdf:Property; object-valued (i.e. rdfs:range must be an Object class) • owl:DatatypeProperty – subclass of rdf:Property; datatype-valued (i.e. its rdfs:range must be an rdfs:Datatype) ⇒ OWL ontologies require each property to be typed in such a way! (for reasons of space sometimes omitted in examples) According to their Cardinality • specifying n:1 or 1:n cardinality: owl:FunctionalProperty, owl:InverseFunctionalProperty ⇒ useful for deriving that objects must be different from each other. According to their Properties

332

# FUNCTIONAL CARDINALITY SPECIFICATION

property rdf:type owl:FunctionalProperty

- not a constraint, but
- if such a property results in two things ... these things are inferred to be the same.

```
@prefix : <foo://bla/names#>.
@prefix persons: <foo://bla/persons/>.
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#> .
@prefix owl: <http://www.w3.org/2002/07/owl#>.
:world :has_pope persons:jorgebergoglio .
:world :has_pope [ :name "Franziskus" ] .
:has_pope rdf:type owl:FunctionalProperty.
[Filename: RDF/popes.n3]
prefix : <foo://bla/names#>
```

```
prefix persons: <foo://bla/persons/>
```

```
select ?N from <file:popes.n3>
```

```
where { persons:jorgebergoglio :name ?N }
```

[Filename: RDF/popes.sparql]

# **OWL: RESTRICTION – EXAMPLE** • owl:Restriction for $\exists p.C$ and $\forall p.C$ . (cf. earlier examples) • Definition of "Parent" as $Parent \equiv Person \sqcap \exists hasChild. \top$ (can be used for conclusions in both directions), • Range axiom as constraint: Parent $\sqsubseteq$ $\forall$ hasChild.Person (use only in the " $\Rightarrow$ " direction) @prefix : <foo://bla#>. @prefix family: <foo://bla/persons/>. @prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>. @prefix owl: <http://www.w3.org/2002/07/owl#>. :Parent owl:equivalentClass [ owl:intersectionOf ( :Person [ a owl:Restriction: owl:onProperty :hasChild; owl:minCardinality 1 ] ) ] . :Parent rdfs:subClassOf [ a owl:Restriction; owl:onProperty :hasChild; owl:allValuesFrom :Person ] . family:john a :Person; :hasChild family:alice . family:sue a :Parent . [Filename: RDF/restriction.n3]

```
334
```

owl:Restriction – Example (cont'd)
prefix : <foo: bla#=""></foo:>
select ?X ?CC ?Y ?C
<pre>from <file:restriction.n3></file:restriction.n3></pre>
where {{?X a : Person; a ?CC} union {?Y : hasChild ?C}} [File: RDF/restriction.sparql]

- How to check whether it knows that Sue has a child?
  - ... only *implicitly* known resources are never contained in SPARQL answers (impedance mismatch between SPARQL and DL).
  - they are only known *inside* the reasoner.
  - for looking inside the reasoner's "private" knowledge, appropriate auxiliary classes have to be defined in the OWL ontology which are then queried by SPARQL (as in many later examples)
- note also the separation of the domain into notions (<foo://bla#>) and instances (<foo://bla/persons/>).

This will not be cleanly done in the subsequent examples because it costs space.

mou DDE/reatriation rdfl	
r	ne: RDF/restriction.rdf]

336

# RESTRICTIONS (AND OTHER CLASS SPECIFICATIONS) AS SEPARATE BLANK NODES

Consider the following (bad) specification:

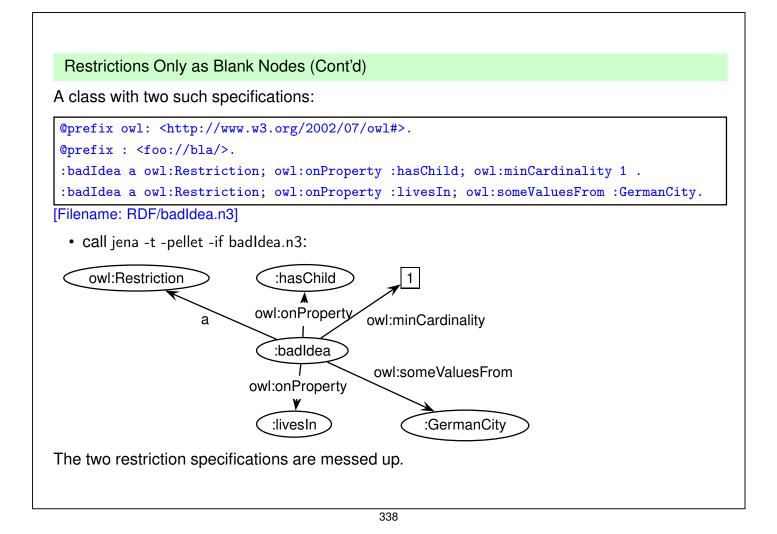
:badIdea a owl:Restriction; owl:onProperty :hasChild; owl:minCardinality 1.

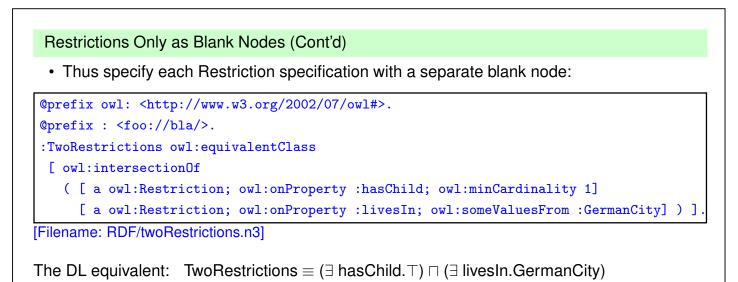
This is not allowed in OWL-DL.

Correct specification:

:badIdea owl:equivalentClass
[a owl:Restriction; owl:onProperty :hasChild; owl:minCardinality 1].

Why? ... there are many reasons, for one of them see next slide.





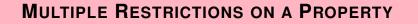
Another reason:

:BadSpecOfParent a owl:Restriction; owl:onProperty :hasChild; owl:minCardinality 1; rdfs:subClassOf :Person.

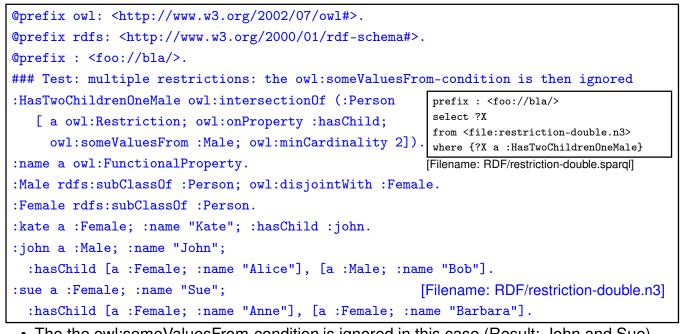
... mixes the *definition* of the Restriction with an assertive axiom:

 $\mathsf{BSOP} \equiv \exists \ge 1 \text{ hasChild}. \top \land \mathsf{ABDE} \sqsubseteq \mathsf{Person}$ 

(This expression probably does not meet the original intention – is *derives* that anything that has a child is made an instance of class "Person"; cf. Slide 329)



- "All persons that have at least two children, and one of them is male"
- first: a straightforward wrong attempt



• The the owl:someValuesFrom-condition is ignored in this case (Result: John and Sue).

340

### Multiple Restrictions on a Property

- "All persons that have at least two children, and one of them is male"
- to expressed as an *intersection* of two separate restrictions:

<pre>@prefix owl: <http: 07="" 2002="" owl#="" www.w3.org="">. @prefix rdfs: <http: 01="" 2000="" rdf-schema#="" www.w3.org="">. @prefix : <foo: bla=""></foo:>.</http:></http:></pre>	<pre>prefix : <foo: bla=""></foo:> select ?X from <file:intersect-restrictions.n3> where {?X a :HasTwoChildrenOneMale}</file:intersect-restrictions.n3></pre>
[ owl:intersectionOf (:Person	[Filename: RDF/intersect-restrictions.sparql]
<pre>[ a owl:Restriction; owl:onProperty :hasChild; owl: [ a owl:Restriction; owl:onProperty :hasChild; owl:</pre>	
<pre>:name a owl:FunctionalProperty. :Male rdfs:subClassOf :Person; owl:disjointWith :Femal</pre>	e.
:Female rdfs:subClassOf :Person. :kate a :Female; :name "Kate"; :hasChild :john.	
<pre>:john a :Male; :name "John";</pre>	e "Bob"].
<pre>:sue a :Female; :name "Sue"; [Filenan :hasChild [a :Female; :name "Anne"], [a :Female; :name</pre>	ne: RDF/intersect-restrictions.n3] me "Barbara"].

 Note: this is different from Qualified Range Restrictions such as "All persons that have at least two male children" – see Slide 402.

### USE OF A DERIVED CLASS

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix : <foo://bla#>.
:kate :name "Kate"; :hasChild :john.
:john :name "John"; :hasChild :alice.
:alice :name "Alice".
:Parent a owl:Class; owl:equivalentClass
 [ a owl:Restriction; owl:onProperty :hasChild; owl:minCardinality 1].
:Grandparent owl:equivalentClass
 [a owl:Restriction; owl:onProperty :hasChild; owl:someValuesFrom :Parent].
[Filename: RDF/grandparent.n3]
prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>
prefix : <foo://bla#>
select ?A ?B
from <file:grandparent.n3>
where {{?A a :Parent} UNION
       {?B a :Grandparent} UNION
       {:Grandparent rdfs:subClassOf :Parent}}
[Filename: RDF/grandparent.spargl]
```

342

### **NON-EXISTENCE OF PROPERTY FILLERS (POSSIBLE SYNTAXES)** @prefix owl: <http://www.w3.org/2002/07/owl#>. @prefix : <foo://bla#>. :ChildlessA owl:intersectionOf (:Person [ a owl:Restriction; owl:onProperty :hasChild; owl:maxCardinality 0]). :ChildlessB owl:intersectionOf (:Person [ a owl:Restriction; owl:onProperty :hasChild; owl:allValuesFrom owl:Nothing]). :ParentA owl:intersectionOf (:Person [owl:complementOf :ChildlessA]). ### (\*) :ParentB owl:intersectionOf (:Person [ a owl:Restriction; owl:onProperty :hasChild; owl:minCardinality 1]). :name a owl:FunctionalProperty. :john a :Person; :name "John"; :hasChild :alice, :bob. :sue a :ParentA; :name "Sue". [Filename: RDF/parents-childless.n3] :george a :Person; a :ChildlessA; :name "George".

- export class tree: ChildlessA and ChildlessB are equivalent,
- · ParentA and ParentB are also equivalent
- note: due to the Open World Assumption, only George is definitely known to be childless.
- Persons where parenthood is not known (Alice, Bob) are neither in Childless nor in Parent!

Note: (\*) states "Parent" vs. "Childless" as a disjoint, total partition of "Person", but it is not *known* to which partition Alice and Bob belong. Both would be possible.

### NON-EXISTENCE OF PROPERTY FILLERS – OPEN WORLD VS. CLOSED WORLD

- · basically the same, Parent and Childless as classes, more persons,
- the focus is now on the different explicit and implicit knowledge about them:

<pre>@prefix owl: <http: 07="" 2002="" owl#="" www.w3.org="">.</http:></pre>
<pre>@prefix : <foo: bla#="">.</foo:></pre>
:Childless owl:intersectionOf (:Person
[ a owl:Restriction; owl:onProperty :hasChild; owl:maxCardinality 0]).
:Parent owl:intersectionOf (:Person
[ a owl:Restriction; owl:onProperty :hasChild; owl:minCardinality 1]).
:name a owl:FunctionalProperty.
<pre>:kate a :Person; :name "Kate"; :hasChild :john, :sue.</pre>
:john a :Person; :name "John"; :hasChild :alice, :bob.
:alice a :Person; :name "Alice".
:bob a :Person; :name "Bob".
:sue a :Parent; :name "Sue".
:george a :Person; a :Childless; :name "George". [Filename: RDF/childless.n3]

344

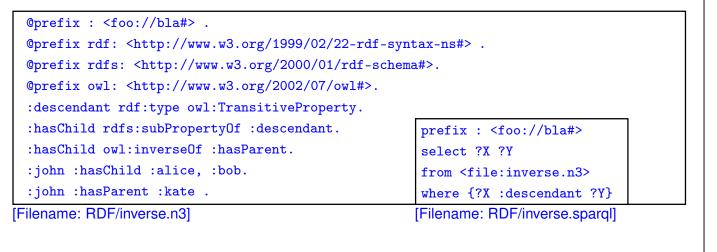
prefix : <foo://bla#>
select ?CL ?NCL ?P ?NP ?NHC ?X ?Y from <file:childless.n3>
where { {?CL a :Childless}
union {?NCL a :Person FILTER NOT EXISTS { ?NCL a :Childless}}
union {?P a :Person FILTER NOT EXISTS { ?NP a :Parent}}
union {?NP a :Person FILTER NOT EXISTS { ?NP a :Parent}}
union {?X :hasChild ?Y}
[Filename: RDF/childless.sparql]
union {?NHC a :Person FILTER NOT EXISTS {?NHC :hasChild ?X}}

DL (and OWL) – everything that is done *inside the reasoner*: open world – monotonic, SPARQL: closed-world – non-monotonic:

- ?CL: only George is known to be Childless.
- ?NCL: Closed-World-Complement of ?C all persons where it cannot be proven that they
  are childless "definitely not childless or maybe not childless" "where it is consistent to
  assume that they are not childless" non-monotonic (all except George).
- Parents ?P: Sue, Kate, John;
- ?NP: Closed-World-Complement of ?P ("consistent to be non-parents" George, Alice, Bob)
- ?X, ?Y: only explicitly known parents/children (Sue not mentioned).
- ?NHC: George, Alice, Bob and Sue(!) no children of them are *explicitly known*.

### **INVERSE PROPERTIES**

- *owl:ObjectProperty* owl:inverseOf *owl:ObjectProperty*
- owl:DatatypeProperties cannot have an inverse (this would define properties of objects, cf. next slide)



346

### No Inverses of owl:DatatypeProperties!

- an owl:DatatypeProperty must not have an inverse:
- ":john :age 35" would imply "35 :ageOf :john" which would mean that a literal has a property, which is not allowed.

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>.
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#> .
@prefix : <foo://bla#> .
# :john :name "John"; :age 35; :hasChild [:name "Alice"], [:name "Bob"; :age 8].
:age a owl:DatatypeProperty.
:hasChild a owl:ObjectProperty.
:parent owl:inverseOf :hasChild.
:ageOf owl:inverseOf :age.
[Filename: RDF/inverseDTProp.n3]
jena -e -pellet -if inverseDTProp.n3
```

```
jena -e -pellet -if inverseDTProp.n3
WARN [main] (OWLLoader.java:352) - Unsupported axiom:
Ignoring inverseOf axiom between foo://bla#ageOf (ObjectProperty)
and foo://bla#age (DatatypeProperty)
```

```
SPECIFICATION OF INVERSE FUNCTIONAL PROPERTIES
 • Mathematics: a mapping m is inverse-functional if the inverse of m is functional:
  x p y is inverse-functional, if for every y, there is at most one x such that xpy holds.
 • Example:
   - hasCarCode is functional: every country has one car code,
   - hasCarCode is also inverse functional: every car code uniquely identifies a country.
 • OWL:
   :m-inverse owl:inverseOf :m .
   :m-inverse a owl:FunctionalProperty .
  not allowed for e.g. mon:carCode a owl:DatatypeProperty:
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo:bla#>.
:carCode a owl:DatatypeProperty; rdfs:domain :Country;
 owl:inverseOf :isCarCodeOf.
# :Germany :carCode "D".
                                                            [Filename: RDF/noinverse.n3]
 • the statement is rejected.
```

```
348
```

### OWL:INVERSEFUNCTIONALPROPERTY

- · such cases are described with owl:InverseFunctionalProperty
- a property P is an owl:InverseFunctionalProperty if  $\forall x, y_1, y_2 : P(y_1, x) \land P(y_2, x) \rightarrow y_1 = y_2$  holds

```
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo:bla#>.
:carCode rdfs:domain :Country; a owl:DatatypeProperty;
    a owl:FunctionalProperty; a owl:InverseFunctionalProperty.
:name a owl:DatatypeProperty; a owl:FunctionalProperty.
:Germany :carCode "D"; :name "Germany".
:DominicanRepublic :carCode "D"; :name "Dominican Republic".
```

```
[Filename: RDF/invfunctional.n3]
```

• the fragment is detected to be inconsistent.

```
OWL:hasKey (OWL 2)
```

Declaration of key attributes  $(k_1, \ldots, k_n)$  is a relevant issue in data modeling.

- a key allows for unambiguously identifying a resource amongst a certain subset of the domain,
- in OWL, keys are not restricted to functional properties (i.e., SQL's UNIQUE is not required),
- values of key properties may be unknown for some instances; they might even be forbidden for some elements of the domain (e.g. using owl:maxCardinality 0 or owl:allValuesFrom owl:Nothing).
- note: InverseFunctionalProperty covers the simple case that n = 1 and the key is global.

```
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo:bla#>.
:name a owl:DatatypeProperty; a owl:FunctionalProperty.
:Country owl:hasKey (:carCode).
:DominicanRepublic a :Country; :carCode "D"; :name "Dominican Republic".
:Germany a :Country; :carCode "D"; :name "Germany". [Filename: RDF/haskey.n3]
```

• the fragment is inconsistent.

350

### OWL:hasKey (OWL 2) for Non-Functional Properties

· keys are not restricted to functional properties:

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo:bla#>.
:District owl:hasKey (:code).
:Country owl:hasKey (:code).
:goettingen a :District; :name "Goettingen"; :code "GOE", "DUD", "HMÃIJ".
:leipzig a :District; :name "Leipzig"; :code "L".
:lahndillkreis a :District; :name "Lahn-Dill-Kreis"; :code "LDK", "DIL", "WZ", "L".
:luxembourg a :Country; :name "Luxembourg"; :code "L".
[Filename: RDF/key-mvd.n3]
```

```
prefix : <foo:bla#>
select ?D ?N ?C
from <file:key-mvd.n3>
where { ?X a ?D ; :name ?N; :code ?C }
```

[Filename: RDF/key-mvd.sparql]

- Lahn-Dill-Kreis and Leipzig are identified (LDK had "L" from 1977-1990).
- Luxembourg is not identified with them since the key definitions are local to districts vs. countries.

OWL:hasKey (OWL 2) for Multi-Property-Keys

- · consider triples about persons found in different Web sources.
- ABSOLUTELY BUGGY (27.7.2017) it equates all four persons below:

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix i <foo:bla#>.
:Person owl:hasKey (:givenName :familyName).
_:b1 a :Person; :givenName "John"; :familyName "Doe"; :age 35 .
_:b2 a :Person; :givenName "John"; :familyName "Doe"; :address "Main Street 1" .
_:b3 a :Person; :givenName "Mary"; :familyName "Doe"; :age 32; :address "Main Street 1" .
_:b4 a :Person; :givenName "Donald"; :familyName "Trump"; :age 70; :address "White House"
#:age a owl:FunctionalProperty.
[Filename: RDF/haskey2.n3]
```

prefix : <foo:bla#>
select ?X ?P ?Y
from <file:haskey2.n3>
where {?X a :Person ; ?P ?Y}

[Filename: RDF/haskey2.sparql]

352

### NAMED AND UNNAMED RESOURCES

(from the DL reasoner's perspective)

Named Resources

- resources with explicit global URIs
   <a href="http://www.semwebtech.org/mondial/10/country/D>">http://www.semwebtech.org/mondial/10/country/D></a>
   <a href="http://bla/bob>">foo://bla/bob></a>
- · resources with local IDs/named blank nodes
- unnamed blank nodes

Unnamed (implicit) Resources

• things that exist only implicitly: John's child in

```
:Parent a owl:Class; owl:equivalentClass
```

```
[ a owl:Restriction; owl:onProperty :hasChild; owl:minCardinality 1].
:john a Parent.
```

• such resources can even have properties (see next slides).

### **Implicit Resources**

- "every person has a father who is a person" and "john is a person".
- the *standard model* is *infinite*: john, john's father, john's father's father, ...
- pure RDF graphs are always finite,
- only with OWL axioms, one can specify such infinite models,
- $\Rightarrow$  they have only finitely many *locally to path length* n different nodes,
  - the reasoner can detect the necessary n ("blocking", cf. Slides 452 ff) and create "typical" different structures.

Aside: "standard model" vs "nonstandard model"

- the term "standard model" is not only "what we understand (in this case)", but is a notion of mathematical theory which –roughly– means "the simplest model of a specification"
- nonstandard models of the above are those where there is a cycle in the ancestors relation.

(as the length of the cycle is arbitrary, this would not make it easier for the reasoner - there is only the possibility to have an owl:sameAs somewhere)

354

### Implicit Resources

@prefix : <foo://bla#>.

@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.

@prefix owl: <http://www.w3.org/2002/07/owl#>.

:Person owl:equivalentClass [a owl:Restriction;

owl:onProperty :father; owl:someValuesFrom :Person].

:bob :name "Bob"; a :Person; :father :john.

:john :name "John"; a :Person.

[Filename: RDF/fathers-and-forefathers.n3]

prefix : <foo://bla#>

prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>

select ?X ?F ?C

from <file:fathers-and-forefathers.n3>

where {{ ?X :father ?F } UNION { ?C a :Person }}

[Filename: RDF/fathers-and-forefathers.sparql]

- Reasoner: works on the model, including blocking, i.e. *modulo equivalence up to paths of length* n.
- SPARQL (and SWRL) rules: works on the graph without the unnamed/implicit resorces.

# 7.3 RDF Graph vs. OWL Model; SPARQL vs. Reasoning

- SPARQL is an RDF (graph) query language
- OWL talks about models.
- Consequences (Overview)
- $\Rightarrow$  SPARQL queries are answered against the graph of triples
  - Some OWL notions are directly represented by triples, such as c a owl:Class.
  - Some others are directly supported by special handling in the reasoners, e.g., *c* rdfs:subClassOf *d* and *c* owl:equivalentClass *d*.
  - some others are only "answered" when given explicitly in the RDF input! The results then do not incorporate further results that could be found by reasoning!
  - OWL notions in the input are often not contained as triples, but are only translated into DL atoms for the reasoner. (e.g. owl:Restriction definitions)
  - Most OWL notions in queries are not "understood" as OWL, but only matched.
  - SPARQL answers are only concerned with the graph, not with implicit things that are only known in the model.

356

# ONTOLOGY LEVEL QUERYING

- SPARQL is defined by *matching* the underlying RDF graph.
- OWL triples are not always part of the RDF graph (they are intended to be translated into DL definitions in the reasoner)
- for traditional DL notions like

```
?C a owl:Class
```

- ?C a rdfs:subClassOf ?D
- ?C owl:equivalentClass ?D
- ?C owl:disjointWith ?D

SPARQL implementations support to translate these internally into DL queries against the reasoner.

• SPARQL-DL (Sirin, Parsia OWLED 2007 [members of the Pellet team]) is a proposal that allows certain further OWL built-ins to be queried.

```
Ontology Level Querying - a practical example
Consider again the "Childless" ontology from Slide 344.
Check that Childless \Box Parent = \emptyset and Person \equiv Childless \sqcup Parent (Partitioning)
  • Allowed: (single line empty bindings result means true)
    prefix : <foo://bla#>
    prefix owl: <http://www.w3.org/2002/07/owl#>
    select ?X from <file:childless.n3>
    where { :Childless owl:disjointWith :Parent } [Filename: RDF/childless1.sparq]
  • Not allowed: complex class expression in the query (empty result since it tries a plain
    match with the RDF data)
    prefix : <foo://bla#>
                                                        [Filename: RDF/childless2.sparql]
    prefix owl: <http://www.w3.org/2002/07/owl#>
                                                       NOT ALLOWED
    select ?X from <file:childless.n3>
    where { :Person owl:equivalentClass [ owl:unionOf (:Childless :Parent) ] }

    instead: add auxiliary class definition to the TBox and export class tree with

    jena -e -if childless.n3 childless3.n3:
    @prefix : <foo://bla#>.
                                                          [Filename: RDF/childless3.n3]
    @prefix owl: <http://www.w3.org/2002/07/owl#>.
```

```
:UnionCLP owl:equivalentClass [ owl:unionOf (:Childless :Parent) ] .
```

# NOT REASONED: OWL:FUNCTIONALPROPERTY

<pre>@prefix rdfs: <http: 01="" 2000="" pre="" rdf-<="" www.w3.org=""></http:></pre>	schema#>.
<pre>@prefix owl: <http: 07="" 2002="" owl#="" www.w3.org=""></http:></pre>	
<pre>@prefix : <foo:bla#>.</foo:bla#></pre>	
:q a owl:FunctionalProperty.	
:p a owl:ObjectProperty; rdfs:domain :D.	
:D owl:equivalentClass [ a owl:Restriction; o	wl:onProperty :p;
o	wl:maxCardinality 1 ].
<pre># :x :p :a, :b. :a owl:differentFrom :b.</pre>	[Filename:RDF/functional.n3]
<pre>prefix owl: <http: 07="" 2002="" owl#="" www.w3.org=""></http:></pre>	
prefix : <foo:bla#></foo:bla#>	
select ?P	
<pre>from <file:functional.n3></file:functional.n3></pre>	
where { ?P a owl:FunctionalProperty }	[Filename:RDF/functional.sparql]
<ul> <li>tries just to match plain { ?P a owl:FunctionalPrope Returns only q.</li> </ul>	erty } triples in the RDF graph.

• does not *derive* that property q is in fact also functiona.

# NOT ALLOWED: COMPLEX TERMS IN SPARQL QUERIES

- example: all cities that are a capital
- works well with pellet alone (June 2017); not allowed with Jena

```
pellet query -query-file countrycaps.sparql \
    mondial-europe.n3 mondial-meta.n3 countrycaps.n3
```

• note: if the answer is empty, check that the mondial-namespace in the used mondial-meta.n3 is correct.

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <http://www.semwebtech.org/mondial/10/meta#> .
:CountryCapital owl:intersectionOf
  (:City [a owl:Restriction; owl:onProperty :isCapitalOf;
        owl:someValuesFrom :Country]). [Filename: RDF/countrycaps.n3]

prefix owl: <http://www.w3.org/2002/07/owl#>
prefix : <http://www.semwebtech.org/mondial/10/meta#>
select ?N1 ?N2
where {{?X a :CountryCapital; :name ?N1} union
        {?Y a [a owl:Restriction; owl:onProperty :isCapitalOf;
        owl:someValuesFrom :Country]; :name ?N2}} [Filename:RDF/countrycaps.sparq0]
```

360

# NOT ALLOWED: COMPLEX TERMS IN SPARQL QUERIES (CONT'D)

- all organizations whose headquarter city is a capital:
- neither allowed by pellet nor by jena+pellet (June 2017; worked with pellet alone in 2013)

```
pellet query -query-file organizations-query2.sparql \
    mondial-europe.n3 mondial-meta.n3
```

### How to do it: Sets of Answers to Queries as Ad-hoc Concepts

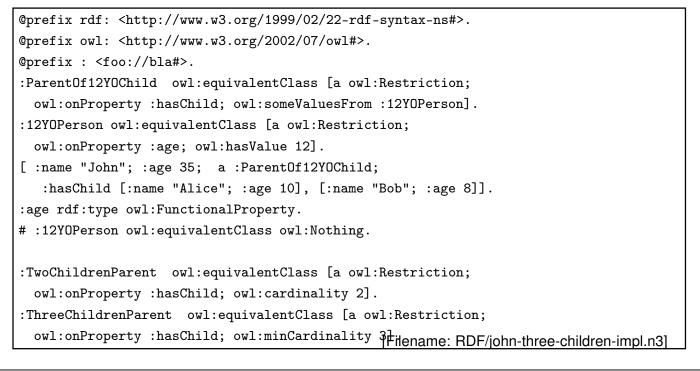
- · The result concept (and maybe others) must be added to the ontology.
- Example: all organizations whose headquarter city is a capital:

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <http://www.semwebtech.org/mondial/10/meta#> .
:CountryCapital owl:equivalentClass
  [ owl:intersectionOf
    (:City [a owl:Restriction; owl:onProperty :isCapitalOf;
           owl:someValuesFrom :Country])].
<bla:Result> owl:equivalentClass [ owl:intersectionOf
  (:Organization [a owl:Restriction; owl:onProperty :hasHeadq;
                                                   [Filename: RDF/organizations-query.n3]
    owl:someValuesFrom :CountryCapital])] .
prefix : <http://www.semwebtech.org/mondial/10/meta#>
select ?A ?N
from <file:organizations-query.n3>
from <file:mondial-europe.n3>
                                                [Filename:RDF/organizations-query.spargl]
from <file:mondial-meta.n3>
where {?X a <bla:Result> . ?X :abbrev ?A . ?X :hasHeadq ?C . ?C :name ?N}
```

```
362
```

### SPARQL ON THE GRAPH: IMPLICITLY KNOWN RESOURCES

 SPARQL does not return any answer related with nodes (=resources) that are only implicitly known (=non-named resources)



### SPARQL and Non-Named Resources (Cont'd)

- implicit resources exist only on the reasoning level,
- not considered by SPARQL queries:

[Filename: RDF/john-three-children-impl.sparql]

- John is a ThreeChildrenParent,
- no person known who is 12 years old
- adding :12YOPerson owl:equivalentClass owl:Nothing makes it inconsistent.
- implicity known things are also not considered for the OWL construct owl:hasKey (cf. Slides 350 and 365) and for SWRL rules (cf. Slides 455 ff).

364

# [ASIDE/EXAMPLE] OWL:HASKEY AND NON-NAMED RESOURCES

Show that owl:hasKey ignores resources that are only implicitly known (OWL ontology see next slide):

- create an (infinite) sequence of implicitly known fathers ... all being persons and having the name "Adam",
- guarantee that the sequence consists of different objects by making it irreflexive. (note: Transitivity and Irreflexivity are not allowed together, thus actually only every person is required to be different from his/her father – the grandfather might be the person again)

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix : <foo:bla#>.
:Person owl:hasKey (:name) .
:name a owl:DatatypeProperty .
# :name a owl:InverseFunctionalProperty . ## that would do it instead of hasKey
:father a owl:FunctionalProperty, owl:IrreflexiveProperty; rdfs:range :Person.
:bob a :Person; :father :john .
:john :name "John" .
:Adam owl:equivalentClass [ a owl:Restriction; owl:onProperty :name; owl:hasValue "Adam"]
:Person rdfs:subClassOf
  [ a owl:Restriction; owl:onProperty :father; owl:someValuesFrom :Adam].
:JohnAdam owl:equivalentClass [ owl:intersectionOf ( :Adam
     [ a owl:Restriction; owl:onProperty :name; owl:hasValue "John"] ) ].
:hasFatherJohnAdam owl:equivalentClass [ a owl:Restriction;
     owl:onProperty :father; owl:someValuesFrom :JohnAdam ] .
:hasGrandpaAdam owl:equivalentClass [ a owl:Restriction; owl:onProperty :father;
    owl:someValuesFrom [ a owl:Restriction; owl:onProperty :father;
       owl:someValuesFrom :Adam ] ].
:AdamFatherAdam owl:equivalentClass [ owl:intersectionOf (:Adam
     [ a owl:Restriction; owl:onProperty :father; owl:someValuesFrom :Adam] ) ] .
[Filename: RDF/forefathers-keys.n3]
```

366

# [ASIDE/EXAMPLE] OWL:HASKEY AND NON-NAMED RESOURCES

```
prefix owl: <http://www.w3.org/2002/07/owl#>
prefix : <foo:bla#>
SELECT ?N ?A ?FA ?AFA ?GPA
FROM <forefathers-keys.n3>
WHERE {{ :bob :father [ :name ?N ] }
    # UNION { ?A :name "Adam" } ## error/bug complains about anon(1)
    UNION { ?FA a :hasFatherJohnAdam }
    UNION { ?AFA a :AdamFatherAdam }
    UNION { ?GPA a :hasGrandpaAdam }}
```

[Filename: RDF/forefathers-keys.sparql]

- implicit nodes are not considered in the answers.
- owl:hasKey is not violated by the fact that several only implicitly known people are named "Adam".

Note that John, being Bob's father, also gets the name "Adam".

### [ASIDE/EXAMPLE] OWL:HASKEY AND NON-NAMED RESOURCES

Another example using multi-attribute keys (which could not be replaced by owl:InverseFunctionalProperty):

- nodes in a (x,y)-coordinate system; consider (10,10)
- insert a pointer to an implicit node (10,10).

368

Aside/Example owl:hasKey and Non-Named Resources (Cont'd)	
prefix owl: <http: 07="" 2002="" owl#="" www.w3.org=""></http:>	
prefix : <foo:bla#></foo:bla#>	
SELECT ?CT ?Y ?T ?SameAsxyxy	
FROM <easykeys-impl.n3></easykeys-impl.n3>	
WHERE {{ :foo :pointTo [ :text ?CT ] }	
UNION { ?Y :text ?T }	
UNION { [:text ?T] }	
UNION { :xyxy owl:sameAs ?SameAsxyxy }}	

- with last in line in source commented out: not much the "pointTo" text is not answered, nothing is :sameAs.
- with last line commented in: the implicit node which is pointed to is equated with :xyxy, made explicit and then equated also with :xy10.

# [ASIDE] OWL VS. RDF LISTS

• RDF provides structures for representing lists by triples (cf. Slide 230): rdf:List, rdf:first, rdf:rest.

These are *distinguished* classes/properties.

- OWL/reasoners have a still unclear relationship with these:
  - use of lists for its internal representation of owl:unionOf, owl:oneOf etc. (that are actually based on collections),
  - do or do not allow the user to query this internal representation,
  - ignore user-defined lists over usual resources.

370

# [ASIDE] UNIONOF (ETC) AS TRIPLES: LISTS

- owl:unionOf (x y z), owl:oneOf (x y z) is actually only syntactic sugar for RDF lists.
- The following are equivalent:

@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>. @prefix owl: <http://www.w3.org/2002/07/owl#>. @prefix : <foo://bla#>. :Male a owl:Class. :Female a owl:Class. :Person a owl:Class; owl:unionOf (:Male :Female). :EqToPerson a owl:Class; owl:unionOf [ a rdf:List; rdf:first :Male; rdf:rest [ a rdf:List; rdf:first :Female; rdf:rest rdf:nil]]. [Filename: RDF/union-list.n3] :x a :Person.

• jena -t -if union-list.n3: both in usual N3 notation as owl:unionOf (:Male :Female).

# [ASIDE] UNIONOF (ETC) AS TRIPLES (CONT'D)

```
prefix owl: <http://www.w3.org/2002/07/owl#>
prefix : <foo://bla#>
select ?C
from <file:union-list.n3>
where {:Person owl:equivalentClass ?C}
```

[Filename: RDF/union-list.sparql]

• jena -q -pellet -qf union-list.sparql: both are equivalent.

```
prefix owl: <http://www.w3.org/2002/07/owl#>
prefix : <foo://bla#>
select ?P1 ?P2 ?X ?Q ?R ?S ?T
from <file:union-list.n3>
where {{:Person owl:equivalentClass :EqToPerson} UNION
    {:Person ?P1 ?X . ?X ?Q ?R . OPTIONAL {?R ?S ?T}} UNION
    {:EqToPerson ?P2 ?X . ?X ?Q ?R} . OPTIONAL {?R ?S ?T}}[Filename: RDF/union-list2.sparq]]
```

 both have actually the same list structure (pellet2/nov 2008: fails; pellet 2.3/sept 2009: fails)

372

# [ASIDE] REASONING OVER LISTS (PITFALLS!)

- rdf:first and rdf:rest are (partially) ignored for reasoning (at least by pellet?); they cannot be used for deriving other properties from it.
- they can even not be used in queries (since pellet2/nov 2008; before it just showed weird behavior)

```
prefix rdf:
```

[Filename: RDF/union-list3.sparql]

- jena-tool with pellet2.3: OK.
- pellet2.3: NullPointerException.

[Aside] Extension of a class defined by a list

Given an RDF list as below, define an owl:Class :Invited which contains exactly the elements in the list (i.e., in the above sample data, :alice, :bob, :carol, :dave).

```
@prefix : <foo:bla#>.
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>.
# Problem: when the real rdf namespace is used, rdf:first/rest are ignored
@prefix rdfL: <http://www.w3.org/1999/02/22-rdf-syntax-nsL#>.
                                                                # <<<<<<<<
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
                                                         prefix : <foo:bla#>
                                                         select ?I
:Invited a owl:Class.
                                                         from <file:invitation-list.n3>
:InvitationList rdfs:subClassOf rdfL:List.
                                                         where {?I a :Invited}
:list1 a :InvitationList; rdfL:first :alice;
  rdfL:rest [a rdfL:List; rdfL:first :bob;
                                                        [Filename: RDF/invitation-list.sparql]
    rdfL:rest [a rdfL:List; rdfL:first :carol;
      rdfL:rest [a rdfL:List; rdfL:first :dave; rdfL:rest rdf:nil]]].
 # rest of an InvitationList is also an InvitationList
:InvitationList owl:equivalentClass
  [a owl:Restriction;
    owl:onProperty rdfL:rest; owl:allValuesFrom :InvitationList],
  [ a owl:Restriction;
    owl:onProperty rdfL:first; owl:allValuesFrom :Invited].
[Filename: RDF/invitation-list.n3]
```

374

# 7.4 Nominals: The O in SHOIQ

### TBox vs. ABox

**Description Logics Terminology** 

Clean separation between TBox and ABox vocabulary:

- TBox: RDFS/OWL vocabulary for information about classes and properties (further partitioned into definitions and axioms),
- ABox: Domain vocabulary and rdf:type.

### RDF/RDF/OWL Ontologies

- Syntactically: allow to mix everything in a single set of triples.
- OWL-DL restriction: clean usage of individuals vs. classes
  - individuals only in application property triples (ABox)
  - classes only in context of RDFS/OWL built-ins (like (X a :Person) or (:hasChild rdfs:range :Person), etc.) (TBox)

**Recall: Reification** 

- Reification treats a class (e.g. :Penguin) or a property as an individual (:Penguin a :Species)
- reification assigns properties from an application domain to classes and properties.
- · useful when talking about metadata notions,
- risk: allows for paradoxes.

### NOMINALS

- use individuals (that usually occur only in the ABox) in *specific positions* in the TBox:
- as individuals (that are often implemented in the reasoner as unary classes) with [a owl:Restriction; owl:onProperty property; owl:hasValue object] (the class of all things such that {?x property object} holds).
- in enumerated classes *class* owl:oneOf (o<sub>1</sub>,...,o<sub>n</sub>) (*class* is defined to be the set {o<sub>1</sub>,...,o<sub>n</sub>}).

```
376
```

### **USING NOMINALS: ITALIAN CITIES**

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix mon: <http://www.semwebtech.org/mondial/10/meta#>.
@prefix it: <foo://bla#>.
it:Italy owl:sameAs <http://www.semwebtech.org/mondial/10/countries/I/>.
it:ItalianCity a owl:Class; owl:intersectionOf
  (mon:City
    [a owl:Restriction; owl:onProperty mon:cityIn;
    owl:hasValue it:Italy]). # Nominal: an individual in a TBox axiom
```

[Filename: RDF/italiancities.n3]

```
prefix it: <foo://bla#>
select ?X ?Y
from <file:mondial-meta.n3>
from <file:mondial-europe.n3>
from <file:italiancities.n3>
where {?X a it:ItalianCity}
```

[Filename: RDF/italiancities.sparql]

 the query {?X :cityIn <http://www.semwebtech.org/mondial/10/countries/l/>} would be shorter, but here a class should be defined for further use ...

# AN ONTOLOGY IN OWL

Consider the Italian-English-Ontology from Slide 52.

- m ·
owl:Thing
bla:Person
bla:English
bla:Hooligan
bla:Gentleman
bla:Italian = bla:Lazy
<pre>owl:Nothing = bla:LatinLover</pre>
<ul> <li>LatinLover is empty,</li> </ul>
thus Italian $\equiv$ Lazy.

378

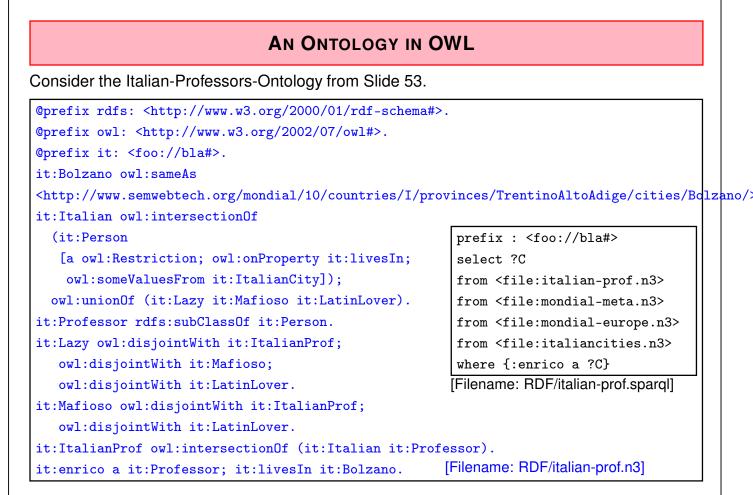
Italians and Englishmen (Cont'd)

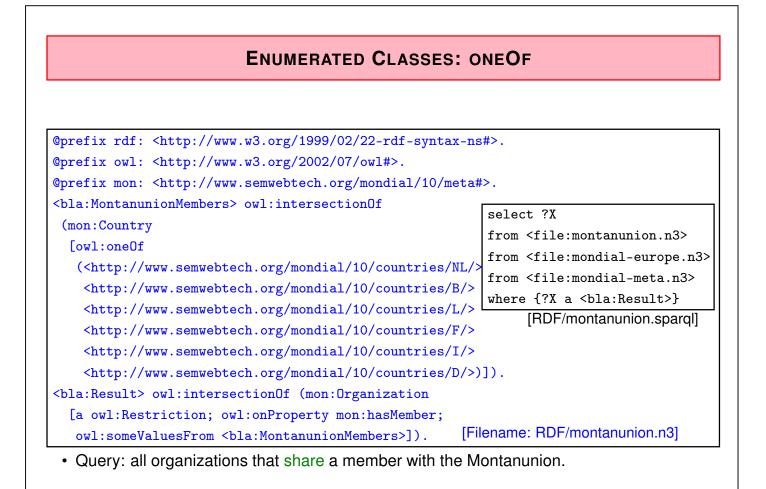
• the conclusions apply to the instance level:

```
@prefix : <foo://bla#>.
:mario a :Italian.
[Filename: RDF/mario.n3]
```

```
prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>
prefix : <foo://bla#>
select ?C
from <file:italian-english.n3>
from <file:mario.n3>
where {:mario rdf:type ?C}
```

[Filename: RDF/italian-english.sparql]





oneOf (Example Cont'd) previous example: "all organizations that share a member with the Montanunion." (DL:  $x \in \exists$ hasMember.MontanunionMembers) • "all organizations where all members are also members of the Montanunion." (DL:  $x \in \forall$ hasMember.MontanunionMembers) The result is empty (although there is e.g. BeNeLux) due to open world: it is not known whether there may exist additional members of e.g. BeNeLux. • Only if the membership of Benelux is "closed", results can be proven: @prefix owl: <http://www.w3.org/2002/07/owl#>. @prefix mon: <http://www.semwebtech.org/mondial/10/meta#>. <http://www.semwebtech.org/mondial/10/organizations/Benelux/> a [a owl:Restriction; select ?X owl:onProperty mon:hasMember; owl:cardinality 3]. from <file:montanunion.n3> <bla:SubsetOfMU> owl:intersectionOf (mon:Organization from <file:montanunion2.n3> [a owl:Restriction; owl:onProperty mon:hasMember; from <file:mondial-europe.n3> owl:allValuesFrom <bla:MontanunionMembers>]). from <file:mondial-meta.n3>

382

mon:name a owl:FunctionalProperty. # not yet given in th where {?X a <bla:SubsetOfMU>}

[RDF/montanunion2.sparql]

### oneOf (Example Cont'd)

[Filename: RDF/montanunion2.n3]

• "all organizations that cover all members of the Montanunion."

```
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix mon: <http://www.semwebtech.org/mondial/10/meta#>.
<bla:EUMembers> owl:equivalentClass [a owl:Restriction;
    owl:onProperty mon:isMember; owl:hasValue
    <http://www.semwebtech.org/mondial/10/organizations/EU/>].
[Filename: RDF/montanunion3.n3]
prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>
select ?X # ?Y ?Z
from <file:montanunion.n3>
from <file:montal-europe.n3>
from <file:mondial-europe.n3>
where {#{?Y a <bla:EUMembers>} UNION {?Z a <bla:MontanunionMembers>} UNION
    {<bla:MontanunionMembers> rdfs:subClassOf ?X}
```

[Filename: RDF/montanunion3.sparql]

# ONEOF (EXAMPLE CONT'D)

Previous example:

- only for one organization
- defined a class that contains all members of the organization
- not possible to define a family of classes one class for each organization.
- this would require a *parameterized constructor*:

" $c_{org}$  is the set of all members of org" Second-Order Logic: each organization can be seen as a unary predicate (=set):  $\forall Org : Org(c) \leftrightarrow hasMember(Org, c)$ or in F-Logic syntax: C isa Org :- Org:organization[hasMember->C] yields e.g.  $I(eu) = \{germany, france, \ldots\}, I(nato) = \{usa, canada, germany, \ldots\}$ Recall that "organization" itself is a predicate:  $I(organization) = \{eu, nato, \ldots, \}$ 

So we have again reification: organizations are both first-order-individuals and classes.

384

# CONVENIENCE CONSTRUCT: OWL:ALLDIFFERENT

- owl:oneOf defines a class as a closed set;
- in owl:oneOf (x<sub>1</sub>, ..., x<sub>n</sub>), two items may be the same (open world),

### owl:AllDifferent

• Triples of the form :a owl:differentFrom :b state that two individuals are different. For a database with n elements, one needs

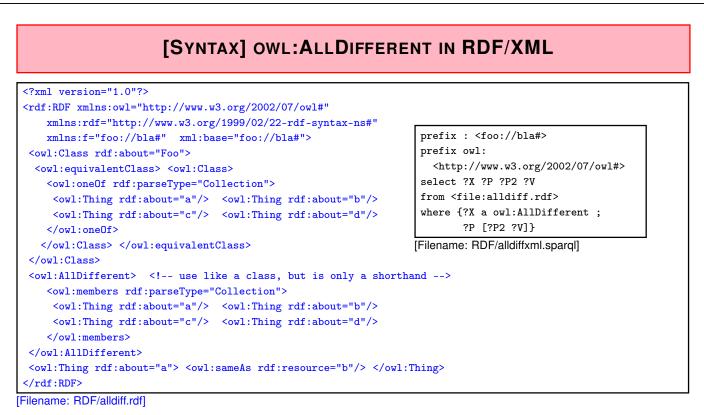
 $(n-1) + (n-2) + \ldots + 2 + 1 = \sum_{i=1..n} i = n \cdot (n+1)/2 = O(n^2)$  such statements.

• The -purely syntactical- convenience construct

```
[ a owl:AllDifferent; owl:members (r_1 r_2 ... r_n) ]
```

provides a shorthand notation.

- it is *immediately* translated into the set of all statements  $\{r_i \text{ owl:differentFrom } r_j \mid i \neq j \in 1..n\}$
- [ a owl:AllDifferent; owl:members (...) ]
   is to be understood as a (blank node) that acts as a *specification* that the listed things are different that does not actually exist in the model.



- AllDifferent is only intended as a kind of command to the application to add all pairwise "different-from" statements, it does not actually introduce itself as triples:
- querying {?X a owl:AllDifferent} is actually not intended.

386

### [SYNTAX] OWL: ALLDIFFERENT IN N3

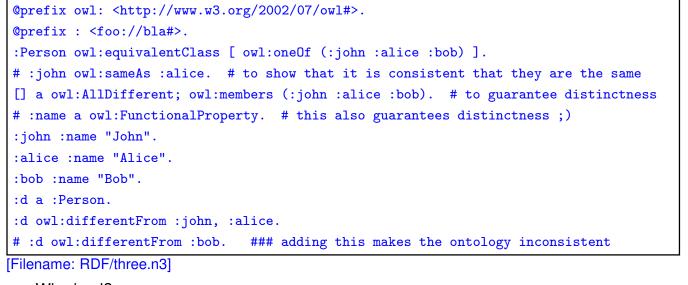
Example:

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo://bla#>.
:Foo owl:equivalentClass [ owl:oneOf (:a :b :c :d) ].
# both the following syntaxes are equivalent and correct:
[ a owl:AllDifferent; owl:members (:a :b)].
[] a owl:AllDifferent; owl:members (:c :d).
:a owl:sameAs :b.
# :b owl:sameAs :d.
[Filename: RDF/alldiff.n3]
```

```
prefix : <foo://bla#>
select ?X ?Y
from <file:alldiff.n3>
where {?X a owl:AllDifferent ; ?P [?P2 ?V]}
[Filename: RDF/alldiff.sparq]]
```

### ONEOF: A TEST

- owl:oneOf defines a "closed set" (use with anonymous class; see below):
- note that in owl:oneOf (x<sub>1</sub>, ..., x<sub>n</sub>), two items may be the same (open world),
- optional owl:AllDifferent to guarantee that  $(x_1, \ldots, x_n)$  are pairwise distinct.



• Who is :d?

388

### oneOf: a Test (cont'd)

Who is :d?

· check the class tree:

bla:Person - (bla:bob, bla:alice, bla:d, bla:john)

The class tree does not indicate which of the "four" identifiers are the same.

• and ask it:

```
prefix : <foo://bla#>
select ?N
from <file:three.n3>
where {:d :name ?N}
[Filename: RDF/three.sparqI]
The answer is ?N/"Bob".
```

# 7.5 Closing Parts of the Open World

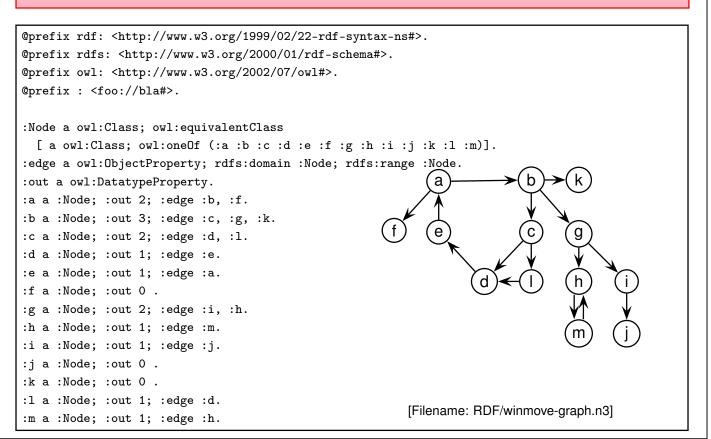
- "forall items" is only applicable if additional items can be excluded (⇒ locally closed predicate/property),
- often, RDF data is generated from a database,
- certain predicates can be closed by defining restriction classes with maxCardinality.

390

### OWL:ALLVALUESFROM

<pre>@prefix owl: <http: 07="" 2002="" owl#="" www.w3.org="">.</http:></pre>
<pre>@prefix : <foo: bla#="">.</foo:></pre>
[ a :Male; a :ThreeChildrenParent; :name "John";
<pre>:hasChild [a :Female; :name "Alice"], [a :Male; :name "Bob"],</pre>
[a :Female; :name "Carol"]].
[ a :Female; a :TwoChildrenParent; :name "Sue";
<pre>:hasChild [a :Female; :name "Anne";], [a :Female; :name "Barbara"]].</pre>
:name a owl:FunctionalProperty.
:OneChildParent owl:equivalentClass [a owl:Re prefix : <foo: bla#=""></foo:>
owl:onProperty :hasChild; owl:cardinality 1 select ?N
:TwoChildrenParent owl:equivalentClass [a owl from <file:allvaluesfrom.n3></file:allvaluesfrom.n3>
owl:onProperty :hasChild; owl:cardinality 2
:ThreeChildrenParent owl:equivalentClass [a of the set rest rest rest rest rest rest rest
<pre>owl:onProperty :hasChild; owl:cardinality 3].</pre>
:OnlyFemaleChildrenParent owl:equivalentClass [a owl:Restriction;
<pre>owl:onProperty :hasChild; owl:allValuesFrom :Female].</pre>
[Filename: RDF/allvaluesfrom.n3]

### EXAMPLE: WIN-MOVE-GAME IN OWL



392

### Win-Move-Game in OWL - the Game Axioms

"If a player cannot move, he loses."

Which nodes are WinNodes, which one are LoseNodes (i.e., the player who has to move wins/loses)?

- if a player can move to some LoseNode (for the other), he will win.
- if a player can move only to WinNodes (for the other), he will lose.
- recall that there can be nodes that are neither WinNodes nor LoseNodes.

```
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>.
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo://bla#>.
"WinNode a owl:Class; owl:intersectionOf ( :Node
    [a owl:Restriction; owl:onProperty :edge; owl:someValuesFrom :LoseNode]).
:LoseNode a owl:Class; owl:intersectionOf ( :Node
    [a owl:Restriction; owl:onProperty :edge; owl:allValuesFrom :WinNode]).
[Filename: RDF/winmove-axioms.n3]
```

### Win-Move-Game in OWL - Closure

```
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>.
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo://bla#>.
:DeadEndNode a owl:Class; rdfs:subClassOf :Node;
  owl:equivalentClass [ a owl:Restriction; owl:onProperty :out; owl:hasValue 0],
                       [ a owl:Restriction; owl:onProperty :edge; owl:cardinality 0].
:OneExitNode a owl:Class; rdfs:subClassOf :Node;
  owl:equivalentClass [ a owl:Restriction; owl:onProperty :out; owl:hasValue 1],
                       [ a owl:Restriction; owl:onProperty :edge; owl:cardinality 1].
:TwoExitsNode a owl:Class; rdfs:subClassOf :Node;
  owl:equivalentClass [ a owl:Restriction; owl:onProperty :out; owl:hasValue 2],
                       [ a owl:Restriction; owl:onProperty :edge; owl:cardinality 2].
:ThreeExitsNode a owl:Class; rdfs:subClassOf :Node;
  owl:equivalentClass [ a owl:Restriction; owl:onProperty :out; owl:hasValue 3],
                       [ a owl:Restriction; owl:onProperty :edge; owl:cardinality 3].
[Filename: RDF/winmove-closure.n3]
```

```
394
```

### Win-Move-Game in OWL: DeadEndNodes

Prove that DeadEndNodes are LoseNodes:

- · obvious: Player cannot move from there
- · exercise: give a formal (Tableau) proof
- The OWL Reasoner does it:

```
prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>
prefix : <foo://bla#>
select ?X
from <file:winmove-axioms.n3>
from <file:winmove-closure.n3>
where {:DeadEndNode rdfs:subClassOf :LoseNode}
[Filename: RDF/deadendnodes.sparql]
```

The answer contains an (empty) tuple which means "yes".